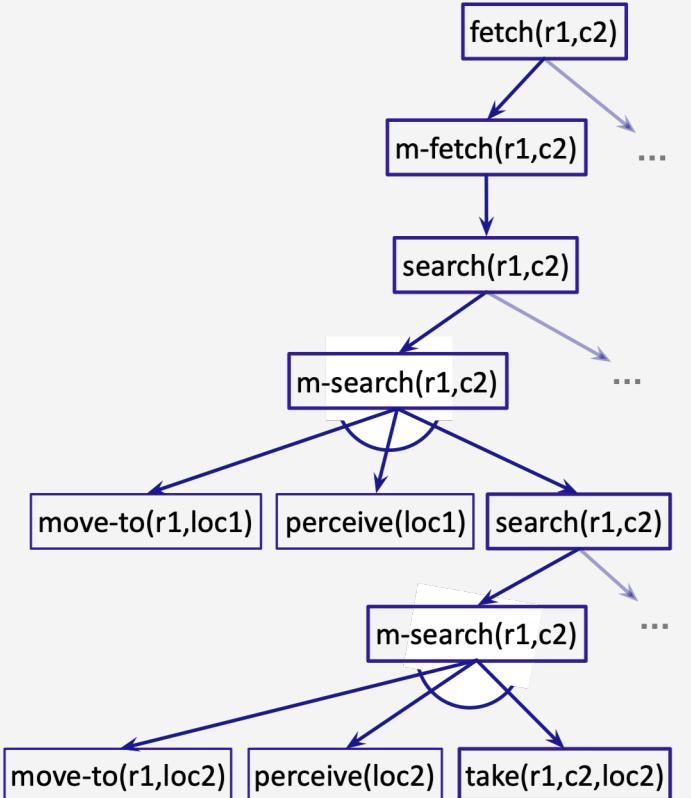


Automated Planning and Acting

Refinement Methods

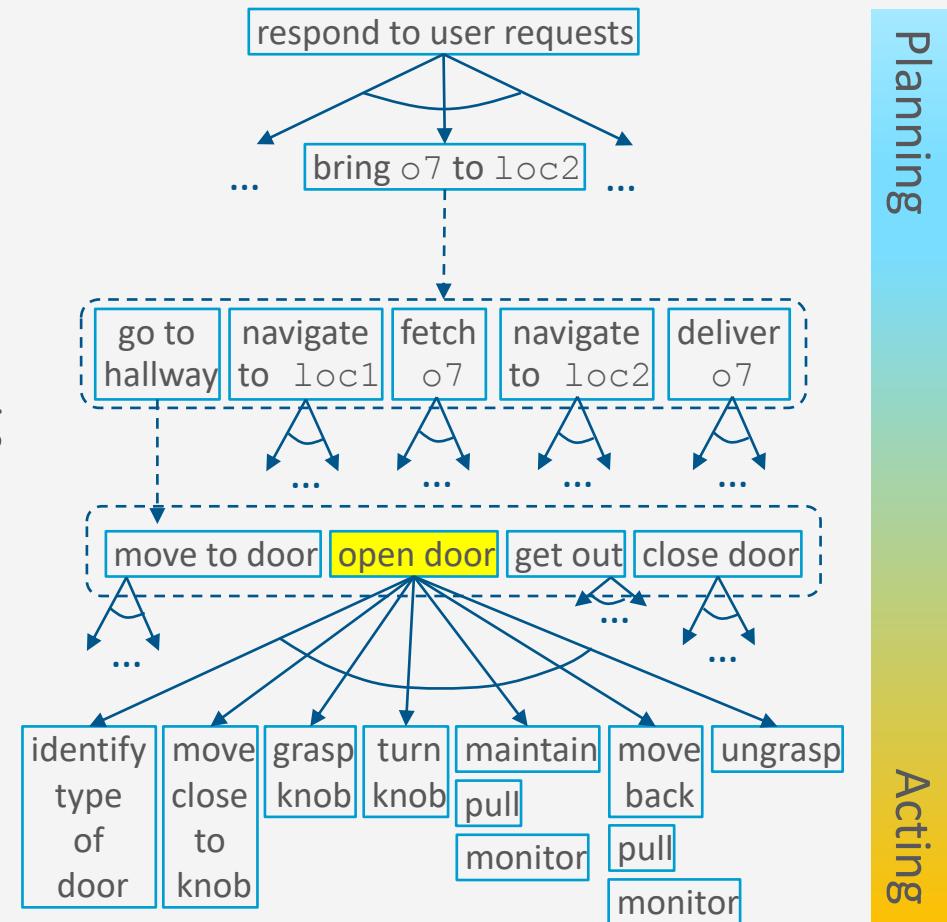


Content: Planning and Acting

1. With Deterministic Models
2. With Refinement Methods
 - Operational Models
 - Refinement-Acting Machine
 - Refinement Planning
 - Acting and Refinement Planning
3. With Temporal Models
4. With Nondeterministic Models
5. With Probabilistic Models
6. By Decision Making
 - A. Foundations
 - B. Extensions
 - C. Structure
7. With human-awareness

Motivation & Assumptions

- Hierarchically organised deliberation
 - At high levels, abstract actions
 - At lower levels, more detail
- Refine abstract actions into ways of carrying them out
 - How?
- Remove / weaken assumptions from classical planning
 - Characteristics
 - Dynamic environment
 - Imperfect information
 - Overlapping actions
 - Nondeterminism
 - Hierarchy
 - Discrete and continuous variables



Outline per the Book

3.1 *Representation*

- State variables, commands, refinement methods
- Example

3.2 *Acting*

- RAE (Refinement Acting Engine)
- Example
- Extensions

3.3 *Planning*

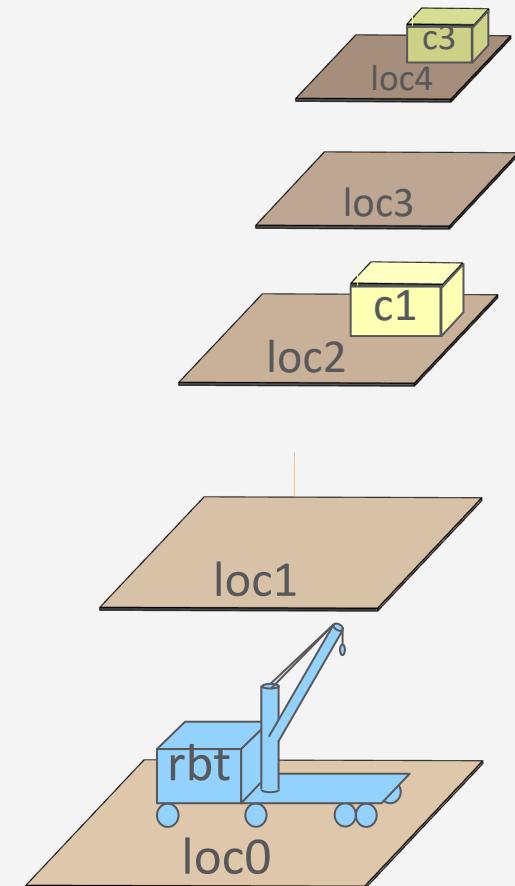
- Motivation and basic ideas
- Deterministic action models
- SeRPE (Sequential Refinement Planning Engine)

3.4 *Using Planning in Acting*

- Techniques
- Caveats

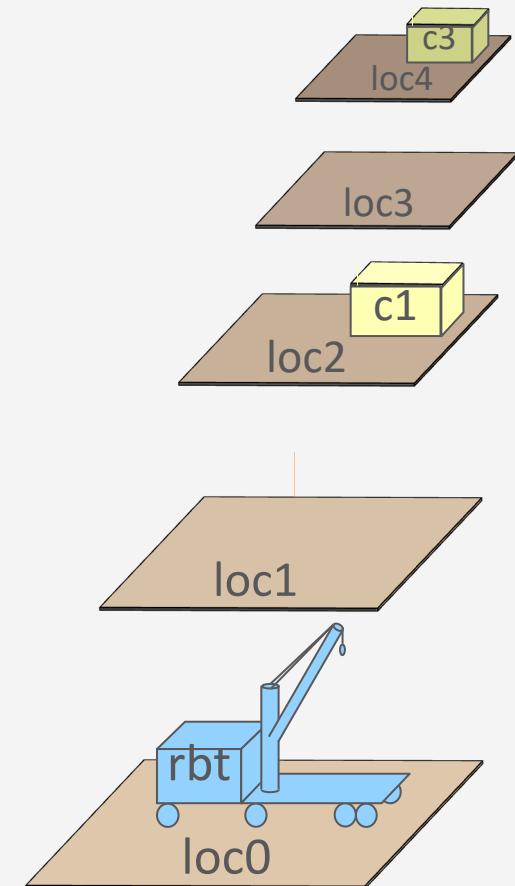
State-variable Representation (Recap)

- Objects:
 - $Robots = \{rbt\}$
 - $Containers = \{c1, c2, c3, \dots\}$
 - $Locations = \{loc0, loc1, loc2, \dots\}$
- State variables: syntactic terms to which we can assign values
 - $loc(r) \in Locations$
 - $load(r) \in Containers \cup \{nil\}$
 - $pos(c) \in Locations \cup Robots \cup \{unknown\}$
 - $view(r, l) \in \{T, F\}$
 - whether robot r has looked at location l
 - r can only see what is at its current location
- State: assign a value to each state variable
 - $\{loc(rbt) = loc0, pos(c1) = loc2, pos(c3) = loc4, pos(c2) = unknown, \dots\}$



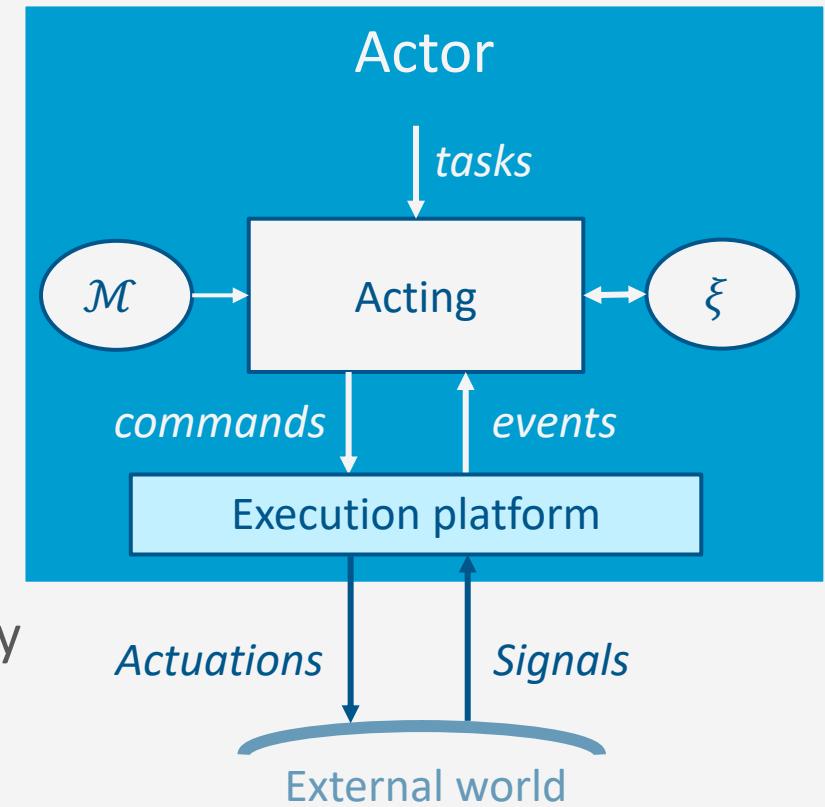
State-variable Representation: Extensions

- Range $\text{ran}(x)$
 - Can be finite, infinite, continuous, discontinuous, vectors, matrices, other data structures
- Assignment statement $x \leftarrow \text{expr}$
 - Expression expr returns a ground value in $\text{ran}(x)$ and has no side-effects on the current state
- Tests (e.g., preconditions)
 - Simple: $x = v, x \neq v, x > v, x < v$
 - Compound: conjunction, disjunction, or negation of simple tests



Commands

- **Command**: primitive function that the execution platform can perform
 - $take(r, o, l)$: robot r takes object o at location l
 - $put(r, o, l)$: r puts o at location l
 - $perceive(r, l)$: robot r perceives what objects are at l
 - r can only perceive what is at its current location
- **Event**: occurrence detected by execution platform
 - $event-name(args)$
 - Exogenous changes in the environment to which the actor may have to react
 - E.g., emergency signal, arrival of transportation vehicle
 - For later: \mathcal{M} : library of methods, ξ : current state (abstraction)

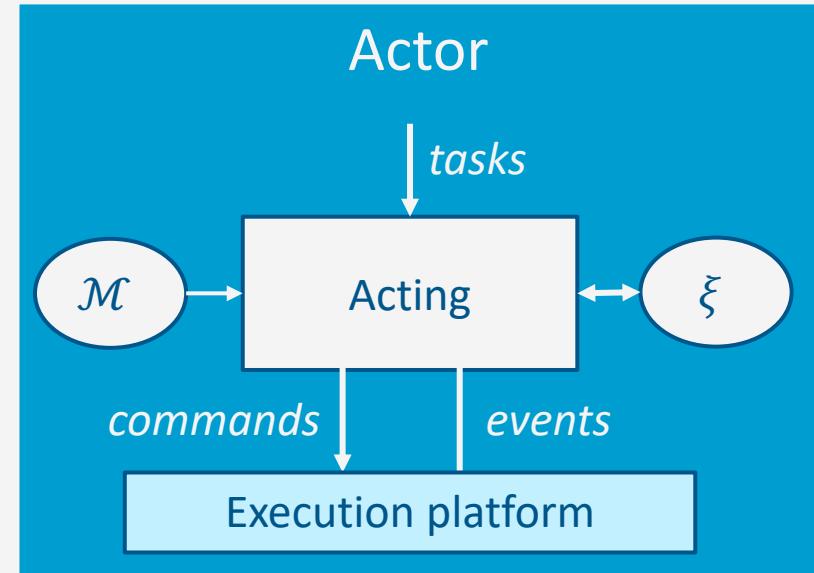


Tasks and Methods

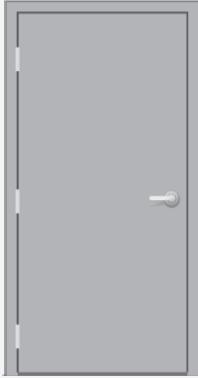
- **Task**: an activity for the actor to perform
 - Could be an abstract action of a plan
- For each task, a set of **refinement methods**
 - **Operational** models:
 - Tell *how* to perform the task
 - Do not predict *what* it will do

```
method-name(arg1, ..., argk)
task: task-identifier
pre: test
body: a program
```

- 
- assignment statements
 - control constructs: if-then-else, while, ...
 - tasks (can extend to include events, goals)
 - commands to the execution platform



Example: “open door” task



- What kind:
 - Hinged on left
 - Opens toward us
 - Lever handle
- Refinement method

```
m-opendoor(r,d,l,h)
task: opendoor(r,d)
pre: loc(r) = l ∧ adj(l,d)
      ∧ handle(d,h)
body:
    while ¬reachable(r,h) do
        move-close(r,h)
        monitor-status(r,d)
        if door-status(d)=closed then
            unLatch(r,d)
            throw-wide(r,d)
        end-monitor-status(r,d)
```

```
m1-unlatch(r,d,l,o)
task: unlatch(r,d)
pre: loc(r,l) ∧ toward-side(l,d) ∧
      side(d,left) ∧ type(d,rotate) ∧ handle(d,o)
body: grasp(r,o)
      turn(r,o,alpha1)
      pull(r,val1)
      if door-status(d)=cracked then ungrasp(r,o)
          else fail

m1-throw-wide(r,d,l,o)
task: throw-wide(r,d)
pre: loc(r,l) ∧ toward-side(l,d) ∧
      side(d,left) ∧ type(d,rotate) ∧
      handle(d,o) ∧ door-status(d)=cracked
body: grasp(r,o)
      pull(r,val1)
      move-by(r,val2)
```

Intermediate Summary

- 3.1 Operational models
 - Tasks, events
 - Commands to the execution platform
 - Extensions to state-variable representation
 - Refinement method: name, task/event, preconditions, body
 - Example: opening a door

Outline per the Book

3.1 *Representation*

- State variables, commands, refinement methods
- Example

3.2 *Acting*

- RAE (Refinement Acting Engine)
- Example
- Extensions

3.3 *Planning*

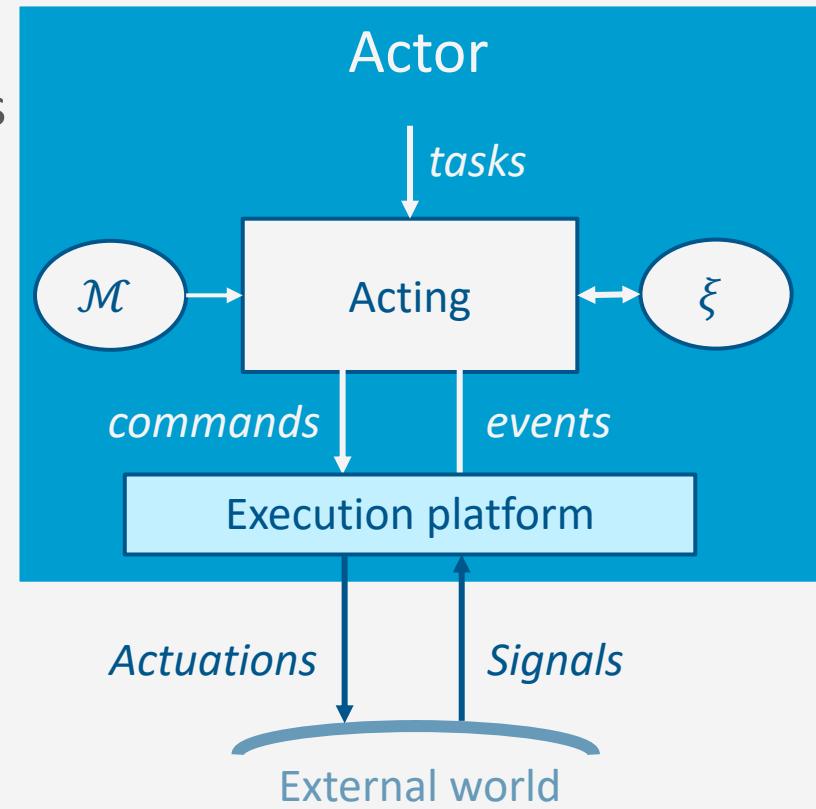
- Motivation and basic ideas
- Deterministic action models
- SeRPE (Sequential Refinement Planning Engine)

3.4 *Using Planning in Acting*

- Techniques
- Caveats

RAE (Refinement Acting Engine)

- Based on OpenPRS programming language
 - Open-source robotics software, deployed in many applications
- Input
 - External tasks, events, current state ξ , library of methods \mathcal{M}
- Output
 - Commands to execution platform
- Perform multiple tasks / events in parallel
 - Purely reactive, no lookahead
- For each task/event, a **refinement stack**
 - Current path in RAE's search tree for the task / event
- **Agenda** = {all current refinement stacks}



RAE (Refinement Acting Engine)

- Input
 - Library of methods \mathcal{M}
 - External tasks and events from an input stream, current state ξ
- Basic idea

loop:

- if new external tasks/events then
 - Add them to Agenda
- for each stack in Agenda
 - Progress it
 - Remove it if it is finished

RAE (\mathcal{M})

$Agenda \leftarrow \emptyset$

loop

```

until the input of external tasks and
events is empty do
  read  $\tau$  in the input stream
   $Candidates \leftarrow Instances(\mathcal{M}, \tau, \xi)$ 
  if  $Candidates = \emptyset$  then
    output("failed to address"  $\tau$ )
  else do
    arbitrarily choose  $m \in Candidates$ 
     $Agenda \leftarrow Agenda \cup \{(\tau, m, \text{nil}, \emptyset)\}$ 
  for each  $stack \in Agenda$  do
    Progress( $stack$ )
    if  $stack = \emptyset$  then
       $Agenda \leftarrow Agenda \setminus \{stack\}$ 

```

Stack element($\tau, m, i, tried$)

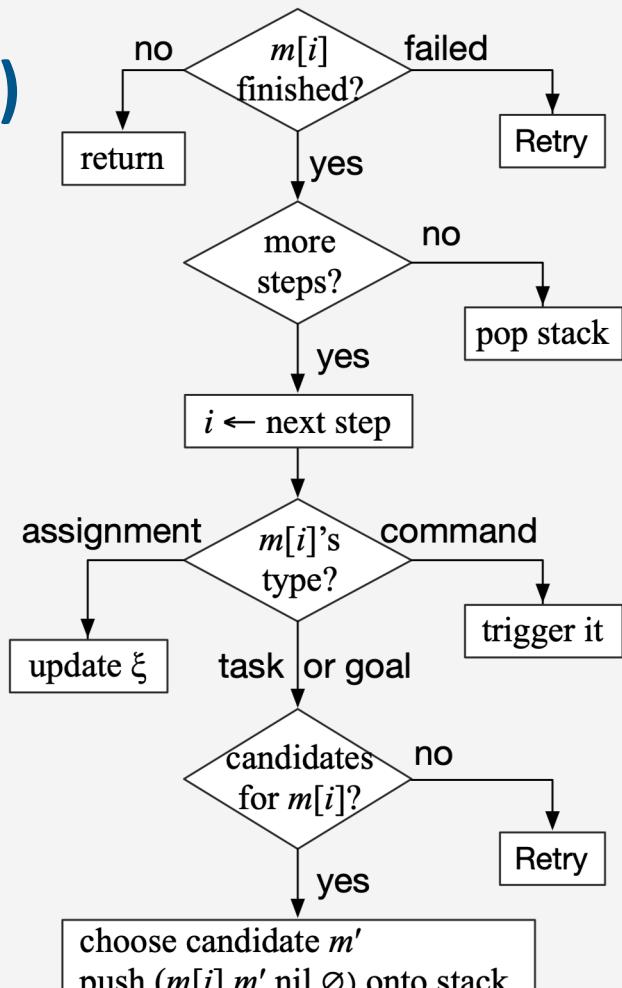
τ task

m instance of a method in \mathcal{M}

i instruction pointer to
 step in body of m

$tried$ method instances already tried

Progress (subroutine)



Just a decision tree

Progress(stack)

```

(τ, m, i, tried) ← top(stack)
if i ≠ nil and m[i] is a command then
    case status(m[i])
        running: return
        failure: Retry(stack); return
        done: continue
    if i is the last step of m then
        pop(stack)
    else do
        i ← nextstep(m, i)
        case type(m[i])
            assignment: update ξ according
                to m[i]; return
            command: trigger m[i]; return
            task or goal: continue
        τ' ← m[i]
        Candidates ← Instances(ℳ, τ', ξ)
        if Candidates = ∅ then
            Retry(stack)
        else do
            arbitrarily choose m' ∈ Candidates
            push((τ', m', nil, ∅), stack)
    
```

Example

m-fetch(r, c)

task: fetch(r, c)

pre:

body:

```
if pos( $c$ ) = unknown then  
    search( $r, c$ )  
else if loc( $r$ ) = pos( $c$ ) then  
    take( $r, c, pos(c)$ )  
else do  
    move-to( $r, pos(c)$ )  
    take( $r, c, pos(c)$ )
```

m-search(r, c)

task: search(r, c)

pre: pos(c) = unknown

body:

```
if  $\exists l$  (view( $r, l$ ) = F) then  
    move-to( $r, l$ )  
    perceive( $l$ )  
    if pos( $c$ ) =  $l$  then  
        take( $r, c, l$ )  
    else search( $r, c$ )  
else fail
```

fetch(r_1, c_2)

?

τ : fetch(r_1, c_2)
 m : ?
 i : (see method)
 $tried: \emptyset$

Refinement stack

Example

m-fetch(r, c)

task: fetch(r, c)

pre:

body:

```
if pos( $c$ ) = unknown then  
    search( $r, c$ )  
else if loc( $r$ ) = pos( $c$ ) then  
    take( $r, c, pos(c)$ )  
else do  
    move-to( $r, pos(c)$ )  
    take( $r, c, pos(c)$ )
```

m-search(r, c)

task: search(r, c)

pre: pos(c) = unknown

body:

```
if  $\exists l$  (view( $r, l$ ) = F) then  
    move-to( $r, l$ )  
    perceive( $l$ )  
    if pos( $c$ ) =  $l$  then  
        take( $r, c, l$ )  
    else search( $r, c$ )  
else fail
```

τ : fetch(r_1, c_2)
 m : m-fetch(r_1, c_2)
 i : (see method)
tried: \emptyset

Refinement stack

fetch(r_1, c_2)

m-fetch(r_1, c_2)

Example

m-fetch(r, c)

task: fetch(r, c)

pre:

body:

if pos(c) = unknown then

search(r, c)

else if loc(r) = pos(c) then

take($r, c, \text{pos}(c)$)

else do

move-to($r, \text{pos}(c)$)

take($r, c, \text{pos}(c)$)

m-search(r, c)

task: search(r, c)

pre: pos(c) = unknown

body:

if $\exists l$ (view(r, l) = F) then

move-to(r, l)

perceive(l)

if pos(c) = l then

take(r, c, l)

else search(r, c)

else fail

τ : search(r_1, c_2)

m : ?

i : (see method)

tried: \emptyset

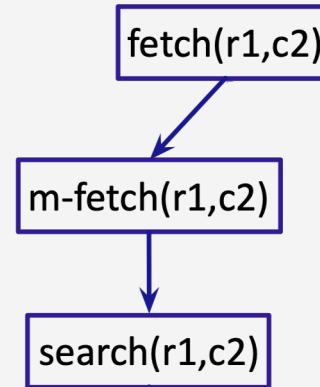
τ : fetch(r_1, c_2)

m : m-fetch(r_1, c_2)

i : (see method)

tried: \emptyset

Refinement stack



Example

$m\text{-fetch}(r,c)$

task: $\text{fetch}(r,c)$

pre:

body:

```

    if  $\text{pos}(c) = \text{unknown}$  then
        search( $r,c$ )
    else if  $\text{loc}(r) = \text{pos}(c)$  then
        take( $r,c,\text{pos}(c)$ )
    else do
        move-to( $r,\text{pos}(c)$ )
        take( $r,c,\text{pos}(c)$ )
    
```

$m\text{-search}(r,c)$

task: $\text{search}(r,c)$

pre: $\text{pos}(c) = \text{unknown}$

body:

```

    if  $\exists l (\text{view}(r,l) = F)$  then
        move-to( $r,l$ )
        perceive( $l$ )
        if  $\text{pos}(c) = l$  then
            take( $r,c,l$ )
        else search( $r,c$ )
    else fail
    
```

$\tau: \text{search}(r_1,c_2)$

$m: \text{m-search}(r_1,c_2)$

$i: (\text{see method})$

$\text{tried: } \emptyset$

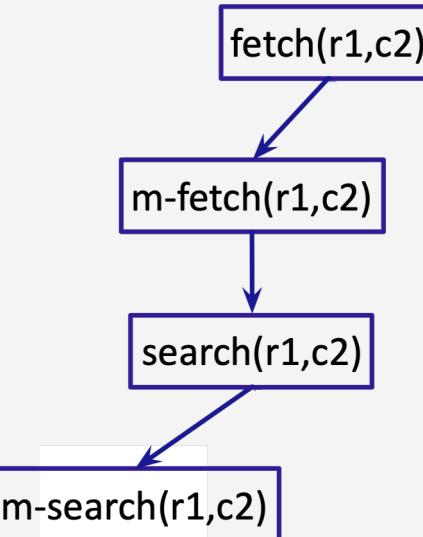
$\tau: \text{fetch}(r_1,c_2)$

$m: \text{m-fetch}(r_1,c_2)$

$i: (\text{see method})$

$\text{tried: } \emptyset$

Refinement stack



Example

m-fetch(r, c)

task: fetch(r, c)

pre:

body:

if $\text{pos}(c) = \text{unknown}$ then

search(r, c)

else if $\text{loc}(r) = \text{pos}(c)$ then

take($r, c, \text{pos}(c)$)

else do

move-to($r, \text{pos}(c)$)

take($r, c, \text{pos}(c)$)

m-search(r, c)

task: search(r, c)

pre: $\text{pos}(c) = \text{unknown}$

body:

if $\exists l (\text{view}(r, l) = F)$ then

move-to(r, l)

perceive(l)

if $\text{pos}(c) = l$ then

take(r, c, l)

else search(r, c)

else fail

...

τ : search(r_1, c_2)

m : m-search(r_1, c_2)

i : (see method)

tried: \emptyset

τ : fetch(r_1, c_2)

m : m-fetch(r_1, c_2)

i : (see method)

tried: \emptyset

Refinement stack

fetch(r_1, c_2)

m-fetch(r_1, c_2)

search(r_1, c_2)

m-search(r_1, c_2)

move-to(r_1, loc_1)

...

...

Example

$m\text{-fetch}(r,c)$

task: $\text{fetch}(r,c)$

pre:

body:

if $\text{pos}(c) = \text{unknown}$ then

search(r,c)

else if $\text{loc}(r) = \text{pos}(c)$ then

take($r,c,\text{pos}(c)$)

else do

move-to($r,\text{pos}(c)$)

take($r,c,\text{pos}(c)$)

$m\text{-search}(r,c)$

task: $\text{search}(r,c)$

pre: $\text{pos}(c) = \text{unknown}$

body:

if $\exists l (\text{view}(r,l) = F)$ then

move-to(r,l)

perceive(l)

if $\text{pos}(c) = l$ then

take(r,c,l)

else search(r,c)

else fail

...

$\tau: \text{search}(r_1,c_2)$

$m: m\text{-search}(r_1,c_2)$

$i: (\text{see method})$

$\text{tried: } \emptyset$

$\tau: \text{fetch}(r_1,c_2)$

$m: m\text{-fetch}(r_1,c_2)$

$i: (\text{see method})$

$\text{tried: } \emptyset$

Refinement stack

fetch(r_1,c_2)

m-fetch(r_1,c_2)

search(r_1,c_2)

m-search(r_1,c_2)

move-to($r_1, \text{loc1}$)

perceive(loc1)

...

Example

$m\text{-fetch}(r,c)$

task: $\text{fetch}(r,c)$

pre:

body:

if $\text{pos}(c) = \text{unknown}$ then

search(r,c)

else if $\text{loc}(r) = \text{pos}(c)$ then

take($r,c,\text{pos}(c)$)

else do

move-to($r,\text{pos}(c)$)

take($r,c,\text{pos}(c)$)

$m\text{-search}(r,c)$

task: $\text{search}(r,c)$

pre: $\text{pos}(c) = \text{unknown}$

body:

if $\exists l (\text{view}(r,l) = F)$ then

move-to(r,l)

~~perceive()~~ *sensor failure*

if $\text{pos}(c) = l$ then

take(r,c,l)

else search(r,c)

else fail

...

$\tau: \text{search}(r_1,c_2)$

$m: m\text{-search}(r_1,c_2)$

$i: (\text{see method})$

$\text{tried: } \emptyset$

$\tau: \text{fetch}(r_1,c_2)$

$m: m\text{-fetch}(r_1,c_2)$

$i: (\text{see method})$

$\text{tried: } \emptyset$

Refinement stack

fetch(r_1,c_2)

m-fetch(r_1,c_2)

search(r_1,c_2)

m-search(r_1,c_2)

move-to($r_1, \text{loc}1$)

perceive($\text{loc}1$)

...

Example

$m\text{-fetch}(r,c)$

task: $\text{fetch}(r,c)$

pre:

body:

if $\text{pos}(c) = \text{unknown}$ then

search(r,c)

else if $\text{loc}(r) = \text{pos}(c)$ then

take($r,c,\text{pos}(c)$)

else do

move-to($r,\text{pos}(c)$)

take($r,c,\text{pos}(c)$)

$m\text{-search}(r,c)$

task: $\text{search}(r,c)$

pre: $\text{pos}(c) = \text{unknown}$

body:

if $\exists l (\text{view}(r,l) = F)$ then

move-to(r,l)

perceive(l)

if $\text{pos}(c) = l$ then

take(r,c,l)

else search(r,c)

else fail

τ : $\text{search}(r_1,c_2)$

m : ?

i : (see method)

$\text{tried}:\{m\text{-search}(r_1,c_2)\}$

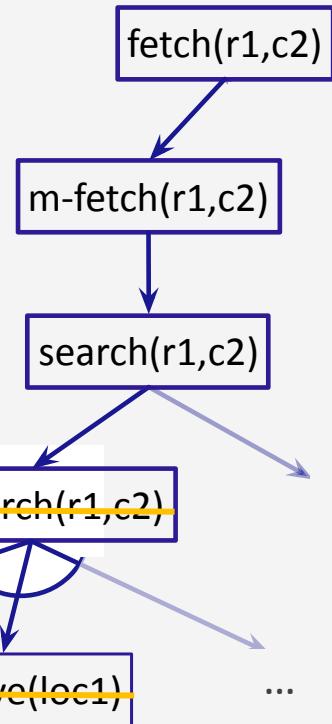
τ : $\text{fetch}(r_1,c_2)$

m : $m\text{-fetch}(r_1,c_2)$

i : (see method)

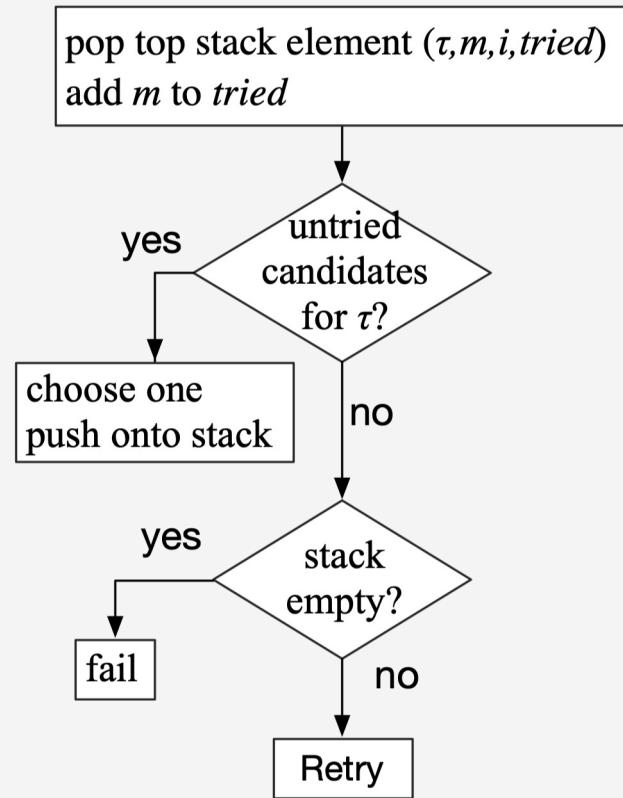
$\text{tried}:\emptyset$

Refinement stack



If other candidates for $\text{search}(r_1,c_2)$, try them.

Retry (subroutine)



Another decision tree

Retry($stack$)

```

 $(\tau, m, i, tried) \leftarrow pop(stack)$ 
 $tried \leftarrow tried \cup \{m\}$ 
 $Candidates \leftarrow Instances(\mathcal{M}, \tau, \xi) \setminus tried$ 
if  $Candidates \neq \emptyset$  then
    arbitrarily choose  $m' \in Candidates$ 
     $push((\tau, m', \text{nil}, \emptyset), stack)$ 
else do
    if  $stack \neq \emptyset$  then
        Retry( $stack$ )
    else do
        output("failed to accomplish"  $\tau$ )
     $Agenda \leftarrow Agenda \setminus stack$ 

```

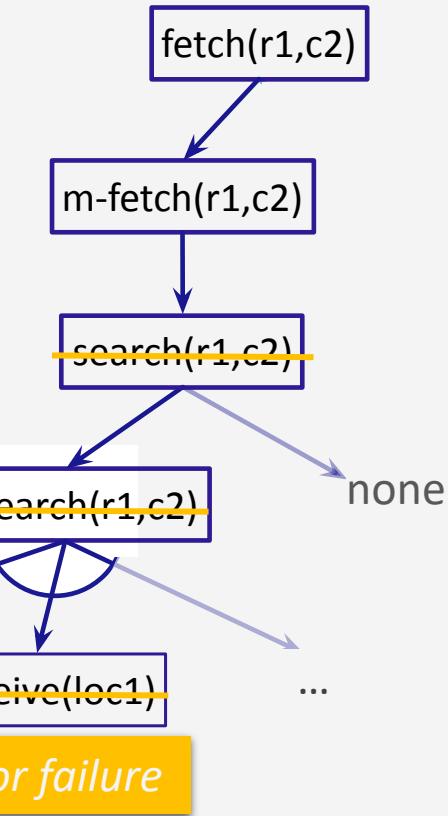
Example

```
m-fetch(r,c)
  task: fetch(r,c)
  pre:
  body:
    if pos(c) = unknown then
      search(r,c)
    else if loc(r) = pos(c) then
      take(r,c,pos(c))
    else do
      move-to(r,pos(c))
      take(r,c,pos(c))
```

```
m-search(r,c)
  task: search(r,c)
  pre: pos(c) = unknown
  body:
    if  $\exists l$  (view(r,l) = F) then
      move-to(r,l)
      perceive(l)
      if pos(c) = l then
        take(r,c,l)
      else search(r,c)
    else fail
```

τ : fetch(r1,c2)
 m : m-fetch(r1,c2)
 i : (see method)
 $tried: \emptyset$

Refinement stack



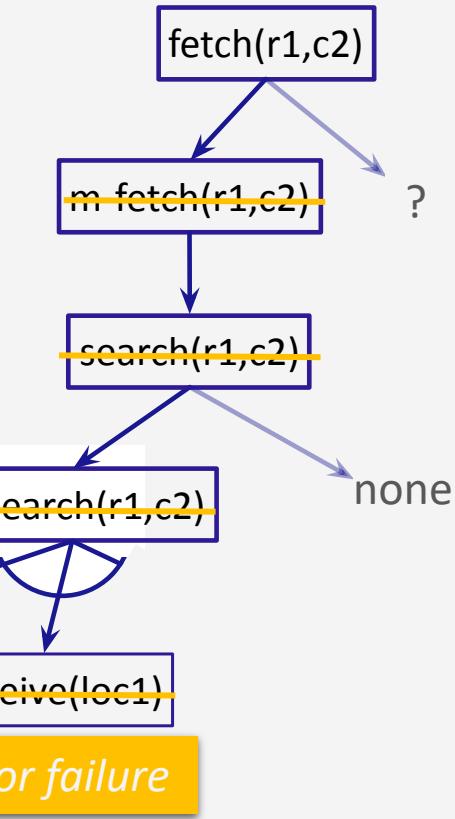
Example

```
m-fetch(r,c)
task: fetch(r,c)
pre:
body:
    if pos(c) = unknown then
        search(r,c)
    else if loc(r) = pos(c) then
        take(r,c,pos(c))
    else do
        move-to(r,pos(c))
        take(r,c,pos(c))
```

```
m-search(r,c)
task: search(r,c)
pre: pos(c) = unknown
body:
    if  $\exists l$  (view(r,l) = F) then
        move-to(r,l)
        perceive(l)
        if pos(c) = l then
            take(r,c,l)
        else search(r,c)
    else fail
```

τ : fetch(r1,c2)
 m : ?
 i : (see method)
tried:{m-fetch(r1,c2)}

Refinement stack



If other candidates for `fetch(r1,c2)`, try them.

Extensions to RAE: Events

- Example: an emergency
 - If r is not already handling another emergency, then
 - Stop what it is doing
 - Go handle the emergency

```
m-emergency(r,l,i)
event: emergency(l,i)
pre:   emergency-handling(r) = F
body:  emergency-handling(r) ← T
       if load(r) ≠ nil then
           put(r,load(r))
           move-to(l)
           address-emergency(l,i)
```

RAE (\mathcal{M})

$Agenda \leftarrow \emptyset$

loop

until the input of external tasks and
events is empty **do**
 read τ in the input stream
 $Candidates \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi)$
if $Candidates = \emptyset$ **then**
 output("failed to address" τ)
else do
 arbitrarily choose $m \in Candidates$
 $Agenda \leftarrow Agenda \cup \{(\tau, m, \text{nil}, \emptyset)\}$
for each $stack \in Agenda$ **do**
 $\text{Progress}(stack)$
if $stack = \emptyset$ **then**
 $Agenda \leftarrow Agenda \setminus \{stack\}$

method-name(arg₁, ..., arg_k)
task: event-identifier
pre: test
body: a program

Extensions to RAE: Goals

- Write as a special task
 - Like others, but includes monitoring: modify Progress
 - if condition becomes true before finishing body(m), stop early
 - if condition isn't true after finishing body(m), fail and try another method

$m\text{-fetch}(r,c)$

task: $\text{fetch}(r,c)$

pre:

body: ~~if $\text{pos}(c) = \text{unknown}$ then
 $\text{search}(r,c)$~~

$\text{achieve}(\text{pos}(c) \neq \text{unknown})$

$\text{move-to}(r, \text{pos}(c))$

$\text{take}(r, c, \text{pos}(c))$

$m\text{-find-where}(r,c)$

task: $\text{achieve}(\text{pos}(c) \neq \text{unknown})$

pre:

body: while exists loc. l s.t. $\text{view}(l) = F$ do

$\text{move-to}(l)$

$\text{perceive}(l)$

RAE (\mathcal{M})

$\text{Agenda} \leftarrow \emptyset$

loop

until the input of external tasks and events is empty **do**

read τ in the input stream

$\text{Candidates} \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi)$

if $\text{Candidates} = \emptyset$ **then**

output("failed to address" τ)

else do

arbitrarily choose $m \in \text{Candidates}$

$\text{Agenda} \leftarrow \text{Agenda} \cup \{(\tau, m, \text{nil}, \emptyset)\}$

for each stack $\in \text{Agenda}$ **do**

Progress(stack)

if stack = \emptyset **then**

$\text{Agenda} \leftarrow \text{Agenda} \setminus \{\text{stack}\}$

method-name(arg₁, ..., arg_k)

task: $\text{achieve}(\text{condition})$

pre: test

body: *a program*

Other Extensions to RAE

- Concurrent subtasks:
 - Refinement stack for each one
- Controlling the progress of tasks:
 - E.g., suspend a task for a while
 - If there are multiple stacks, which ones get higher priority?
 - Application-specific heuristics
- For a task τ , which candidate to try first?
 - Refinement planning

body of a method

...
{concurrent: $\tau_1, \tau_2, \dots, \tau_n$ }

...

$Agenda = \{stack_1, stack_2, \dots, stack_n\}$

$Candidates \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi)$

Intermediate Summary

- 3.2 Refinement Acting Engine (RAE)
 - Purely reactive: select a method and apply it
 - RAE: input stream, *Candidates*, Instances, *Agenda*, refinement stacks
 - Progress: command status, nextstep, type of step
 - Retry: *Candidates \ tried*, *Agenda \ stack*
 - Refinement trees
 - Concurrent tasks: for each, a refinement stack
 - Goal: achieve(condition), uses monitoring
 - Controlling progress, heuristics

Outline per the Book

3.1 *Representation*

- State variables, commands, refinement methods
- Example

3.2 *Acting*

- RAE (Refinement Acting Engine)
- Example
- Extensions

3.3 *Planning*

- Motivation and basic ideas
- Deterministic action models
- SeRPE (Sequential Refinement Planning Engine)

3.4 *Using Planning in Acting*

- Techniques
- Caveats

Motivation

- When dealing with an event or task, RAE may need to make either/or choices
 - *Agenda*: tasks $\tau_1, \tau_2, \dots, \tau_n$
 - Several tasks/events, how to prioritise?
 - Candidates for τ_1 : m_1, m_2, \dots
 - Several candidate methods or commands, which one to try first?
- RAE immediately executes commands
 - Bad choices may be **costly** or **irreversible**

Refinement Planning

- Basic idea:
 - Go step by step through RAE, but do not send commands to execution platform
 - For each command, use a descriptive action model to predict the next state
 - Tells *what*, not *how*
 - Whenever we need to choose a method
 - Try various possible choices, explore consequences, choose best
- Generalisation of Hierarchical Task Network (HTN) planning
 - HTN planning: body of a method is a list of tasks
 - Here: body of method is the same program RAE uses
 - Use it to *generate* a list of tasks

SeRPE (Sequential Refinement Planning Engine)

- SeRPE inputs

\mathcal{M} = {methods}

\mathcal{A} = {action models}

s = initial state

τ = task or goal

- Which candidate method for τ ?

- SeRPE: Nondeterministic choice

- Backtracking point

- How to implement?

- Hierarchical adaptation of backtracking, A*, GBFS, ...

- RAE: Arbitrary choice

- No search, purely reactive

SeRPE ($\mathcal{M}, \mathcal{A}, s, \tau$)

Candidates \leftarrow Instances (\mathcal{M}, τ, s)

if *Candidates* = \emptyset **then**

return failure

 nondeterministically choose $m \in \text{Candidates}$

return Progress-to-finish ($\mathcal{M}, \mathcal{A}, s, \tau, m$)

RAE (\mathcal{M})

Agenda $\leftarrow \emptyset$

loop

until the input of external tasks and events is empty **do**

 read τ in the input stream

Candidates \leftarrow Instances (\mathcal{M}, τ, ξ)

if *Candidates* = \emptyset **then**

 output ("failed to address" τ)

else do

 arbitrarily choose $m \in \text{Candidates}$

Agenda $\leftarrow \text{Agenda} \cup \{(\tau, m, \text{nil}, \emptyset)\}$

for each *stack* $\in \text{Agenda}$ **do**

 Progress (*stack*)

if *stack* = \emptyset **then**

Agenda $\leftarrow \text{Agenda} \setminus \{\text{stack}\}$

SeRPE (Sequential Refinement Planning Engine)

- SeRPE
 - One external task
 - Simulate progressing it all the way to the end
- RAE
 - Several external tasks
 - Each time through loop, progress each one by one step

SeRPE ($\mathcal{M}, \mathcal{A}, s, \tau$)

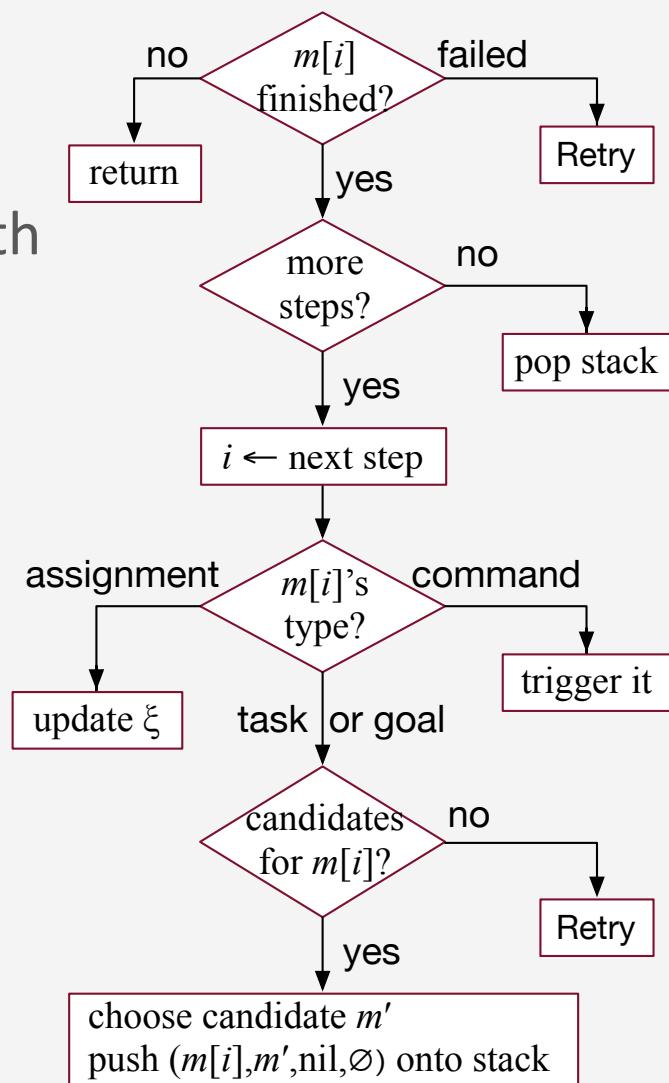
```
Candidates ← Instances ( $\mathcal{M}, \tau, s$ )
if Candidates =  $\emptyset$  then
    return failure
nondeterministically choose  $m \in$  Candidates
return Progress-to-finish ( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
```

RAE (\mathcal{M})

```
Agenda ←  $\emptyset$ 
loop
    until the input of external tasks and
        events is empty do
        read  $\tau$  in the input stream
        Candidates ← Instances ( $\mathcal{M}, \tau, \xi$ )
        if Candidates =  $\emptyset$  then
            output ("failed to address"  $\tau$ )
        else do
            arbitrarily choose  $m \in$  Candidates
            Agenda ← Agenda  $\cup$  { $(\tau, m, \text{nil}, \emptyset)$ }
        for each stack  $\in$  Agenda do
            Progress (stack)
            if stack =  $\emptyset$  then
                Agenda ← Agenda  $\setminus$  {stack}
```

Progress-to-finish

- Like RAE progress with a loop around it
 - Simulates the commands



```

Progress-to-finish( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
   $i \leftarrow \text{nil}; \pi \leftarrow \langle \rangle$ 
  loop
    if  $\tau$  is a goal and  $s \models \tau$  then
      return  $\pi$ 
    if  $i$  is the last step of  $m$  then
      if  $\tau$  is a goal and  $s \not\models \tau$  then
        return failure
      return  $\pi$ 
     $i \leftarrow \text{nextstep}(m, i)$ 
    case type( $m[i]$ )
      assignment:
        update  $s$  according to  $m[i]$ 
      command:
         $a \leftarrow \text{descriptive model of } m[i] \text{ in } \mathcal{A}$ 
        if  $s \models \text{pre}(a)$  then
           $s \leftarrow \gamma(s, a); \pi \leftarrow \pi . a$ 
        else
          return failure
      task or goal:
         $\pi' \leftarrow \text{SeRPE}(\mathcal{M}, \mathcal{A}, s, m[i])$ 
        if  $\pi' = \text{failure}$  then
          return failure
         $s \leftarrow \gamma(s, \pi'); \pi \leftarrow \pi . \pi'$ 
    
```

Progress-to-finish

- Inputs: \mathcal{M} = {methods}, \mathcal{A} = {action models},
 s = initial state, τ = task or goal,
 m = chosen method
- Simulate RAE's goal monitoring
- If $m[i]$ is a command:
 Use action model to predict outcome
- If current step is a task: Call SeRPE recursively
 - Recursion stack \approx RAE's refinement stack
- For failures, no Retry (RAE)
 - A failure means SeRPE could not find a solution
 - Implementation: hierarchical adaptations of backtracking, A*, GBFS, ...

```

Progress-to-finish( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
   $i \leftarrow \text{nil}; \pi \leftarrow \langle \rangle$ 
  loop
    if  $\tau$  is a goal and  $s \models \tau$  then
      return  $\pi$ 
    if  $i$  is the last step of  $m$  then
      if  $\tau$  is a goal and  $s \not\models \tau$  then
        return failure
      return  $\pi$ 
     $i \leftarrow \text{nextstep}(m, i)$ 
    case type( $m[i]$ )
      assignment:
        update  $s$  according to  $m[i]$ 
      command:
         $a \leftarrow$  descriptive model of  $m[i]$  in  $\mathcal{A}$ 
        if  $s \models \text{pre}(a)$  then
           $s \leftarrow \gamma(s, a); \pi \leftarrow \pi.a$ 
        else
          return failure
      task or goal:
         $\pi' \leftarrow \text{SeRPE}(\mathcal{M}, \mathcal{A}, s, m[i])$ 
        if  $\pi' = \text{failure}$  then
          return failure
         $s \leftarrow \gamma(s, \pi'); \pi \leftarrow \pi. \pi'$ 
  
```

Example

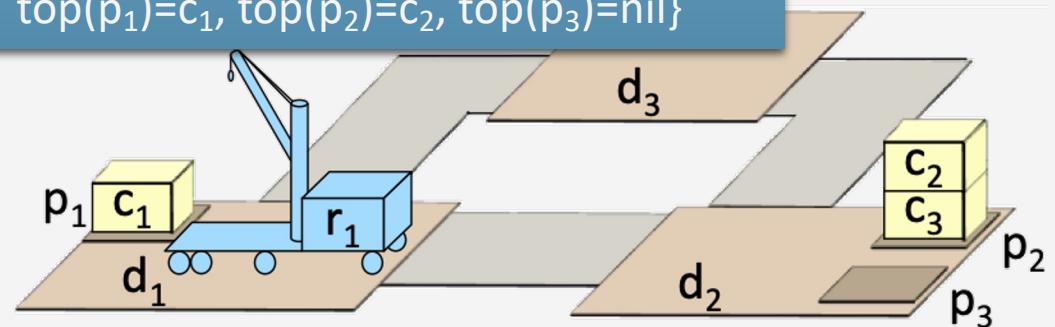
*Candidates = {m1-put-in-pile(c₁,p₂),
m2-put-in-pile(r,c₁,p₁,d,p',d')}*

```
SeRPE( $\mathcal{M}, \mathcal{A}, s, \tau$ )
  Candidates  $\leftarrow$  Instances( $\mathcal{M}, \tau, s$ )
  if Candidates =  $\emptyset$  then
    return failure
  nondeterministically choose  $m \in$  Candidates
  return Progress-to-finish( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
```

m1-put-in-pile(c, p')
task: put-in-pile(c, p')
pre: pile(c)= p'
body: // empty

m2-put-in-pile(r, c, p, d, p', d')
task: put-in-pile(c, p')
pre: pile(c)= p \wedge at(p, d) \wedge at(p', d')
 $\wedge p \neq p' \wedge$ cargo(r)=nil
body: if loc(r) $\neq d$ then navigate(r, d)
uncover(c)
load($r, c, pos(c), p, d$)
if loc(r) $\neq d'$ then navigate(r, d')
unload($r, c, top(p'), p', d$)

$s_0 = \{\text{loc}(r_1)=d_1, \text{cargo}(r_1)=\text{nil}, \text{occupied}(d_1)=\text{T},$
 $\text{occupied}(d_2)=\text{F}, \text{occupied}(d_3)=\text{F},$
 $\text{pos}(c_1)=\text{nil}, \text{pos}(c_2)=c_3, \text{pos}(c_3)=\text{nil},$
 $\text{pile}(c_1)=p_1, \text{pile}(c_2)=p_2, \text{pile}(c_3)=p_2,$
 $\text{top}(p_1)=c_1, \text{top}(p_2)=c_2, \text{top}(p_3)=\text{nil}\}$



Example

task
 $\text{put-in-pile}(c_1, p_2)$
 |
 method
 $\text{m2-put-in-pile}(r_1, c_1, p_1, d_1, p_2, d_2)$

Refinement tree

- The SeRPE pseudocode does not return this, but can easily be modified to do so

SeRPE($\mathcal{M}, \mathcal{A}, s, \tau$)

```

Candidates  $\leftarrow$  Instances( $\mathcal{M}, \tau, s$ )
if Candidates =  $\emptyset$  then
  return failure
nondeterministically choose  $m \in$  Candidates
return Progress-to-finish( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
  
```

$r_1, c_1, p_1, d_1, p_2, d_2$

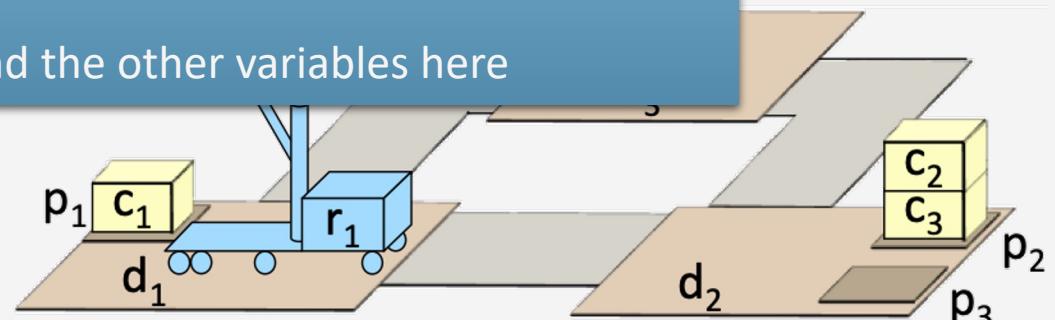
$\text{m2-put-in-pile}(r, c, p, d, p', d')$

task: $\text{put-in-pile}(c, p') \leftarrow$

pre: $\text{pile}(c)=p \wedge \text{at}(p, d) \wedge \text{at}(p', d') \leftarrow$
 $\wedge p \neq p' \wedge \text{cargo}(r)=\text{nil}$

body: if $\text{loc}(r) \neq d$ then $\text{navigate}(r, d)$
 $\text{uncover}(c)$
 $\text{load}(r, c, \text{pos}(c), p, d)$
 $\text{if } \text{loc}(r) \neq d' \text{ then } \text{navigate}(r, d')$
 $\text{unload}(r, c, \text{top}(p'), p', d)$

- m2-put-in-pile starts with $c=c_1$, $p'=p_2$, and r, d, p', d' unbound
- Bind the other variables here



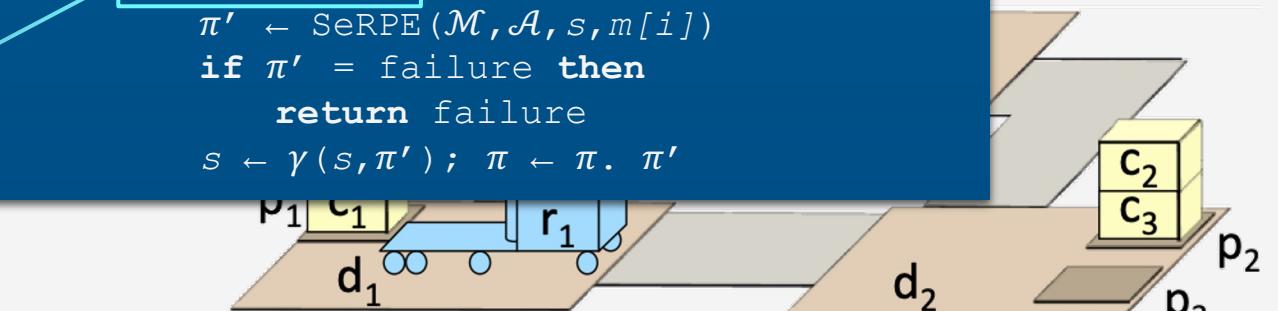
Example

task
 put-in-pile(c_1, p_2)
 |
method
 m2-put-in-pile($r_1, c_1, p_1, d_1, p_2, d_2$)

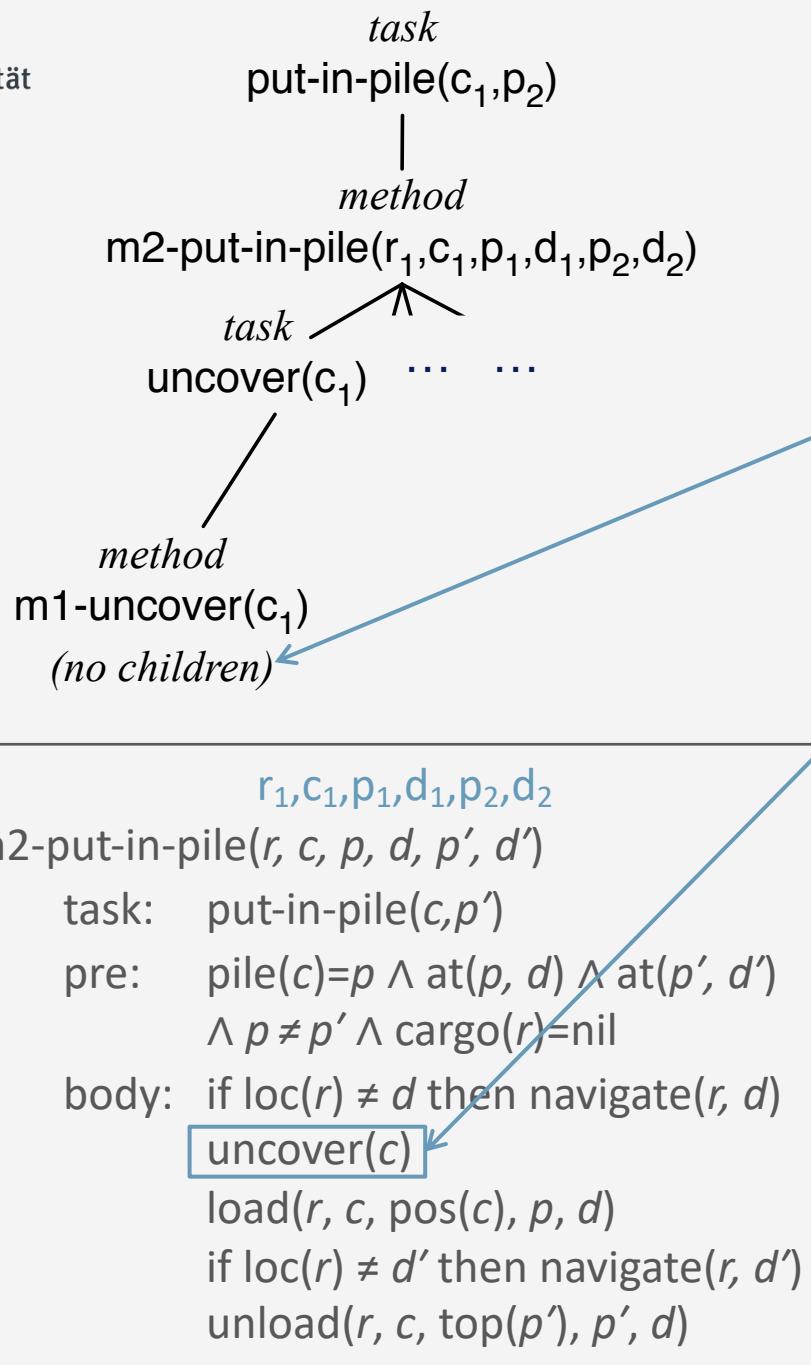
$r_1, c_1, p_1, d_1, p_2, d_2$
 m2-put-in-pile(r, c, p, d, p', d')
 task: put-in-pile(c, p')
 pre: pile(c) = $p \wedge \text{at}(p, d) \wedge \text{at}(p', d')$
 $\wedge p \neq p' \wedge \text{cargo}(r) = \text{nil}$
 body: if $\text{loc}(r) \neq d$ then $\text{navigate}(r, d)$
uncover(c)
 $\text{load}(r, c, \text{pos}(c), p, d)$
 if $\text{loc}(r) \neq d'$ then $\text{navigate}(r, d')$
 $\text{unload}(r, c, \text{top}(p'), p', d)$

```

Progress-to-finish( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
   $i \leftarrow \text{nil}; \pi \leftarrow \langle \rangle$ 
  loop
    if  $\tau$  is a goal and  $s \models \tau$  then
      return  $\pi$ 
    if  $i$  is the last step of  $m$  then
      if  $\tau$  is a goal and  $s \not\models \tau$  then
        return failure
      return  $\pi$ 
     $i \leftarrow \text{nextstep}(m, i)$ 
    case type( $m[i]$ )
      assignment:
        update  $s$  according to  $m[i]$ 
      command:
         $a \leftarrow$  descriptive model of  $m[i]$  in  $\mathcal{A}$ 
        if  $s \models \text{pre}(a)$  then
           $s \leftarrow \gamma(s, a); \pi \leftarrow \pi.a$ 
        else
          return failure
      task or goal:
         $\pi' \leftarrow \text{SeRPE}(\mathcal{M}, \mathcal{A}, s, m[i])$ 
        if  $\pi' = \text{failure}$  then
          return failure
         $s \leftarrow \gamma(s, \pi'); \pi \leftarrow \pi. \pi'$ 
  
```



Example



m1-uncover(c)

task: *uncover(c)*

pre: *top(pile(c))=c*

body: // empty

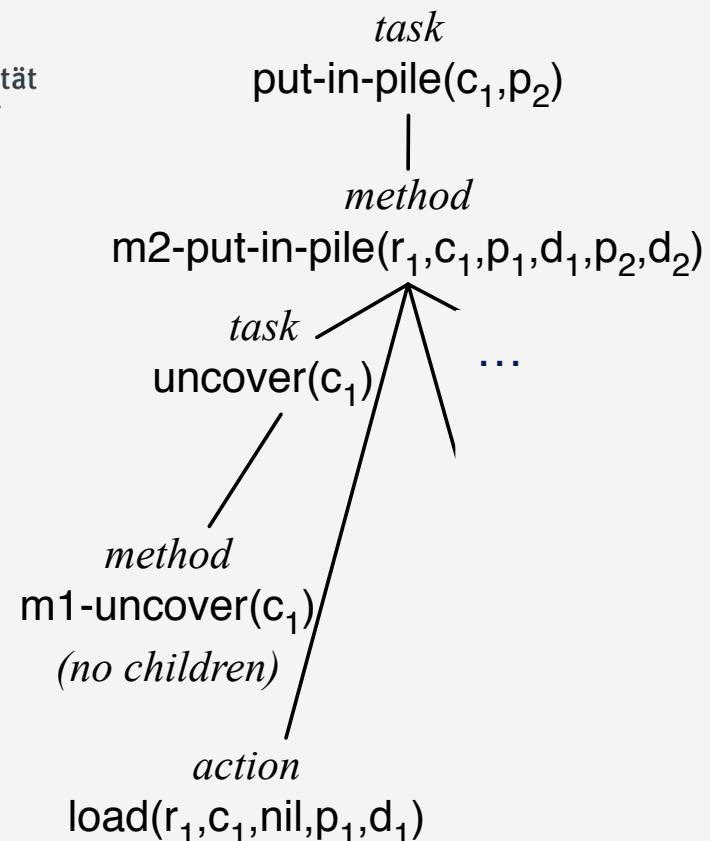
m2-uncover(r, c, p, d)

task: *uncover(c)*

pre: *pile(c)=p \wedge top(p)\neq c*
\wedge at(p, d) \wedge at(p', d) \wedge p'\neq p
\wedge loc(r)=d \wedge cargo(r)=nil

body: while *top(p) \neq c* do
c' \leftarrow top(p)
load(r, c', pos(c'), p, d)
unload(r, c', top(p'), p', d)

Example

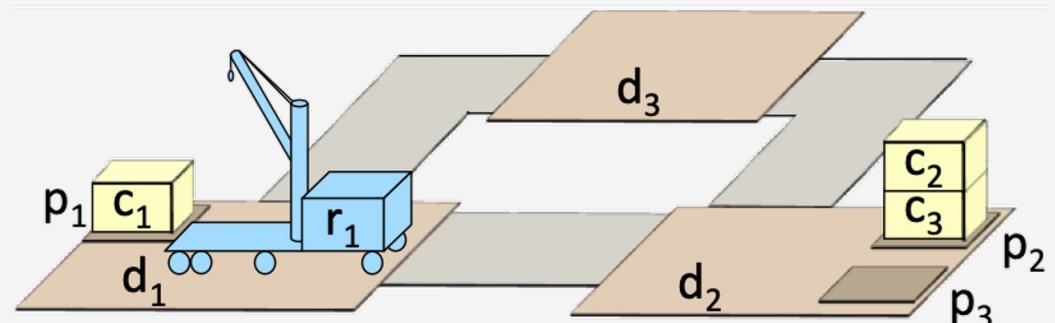


r₁, c₁, p₁, d₁, p₂, d₂

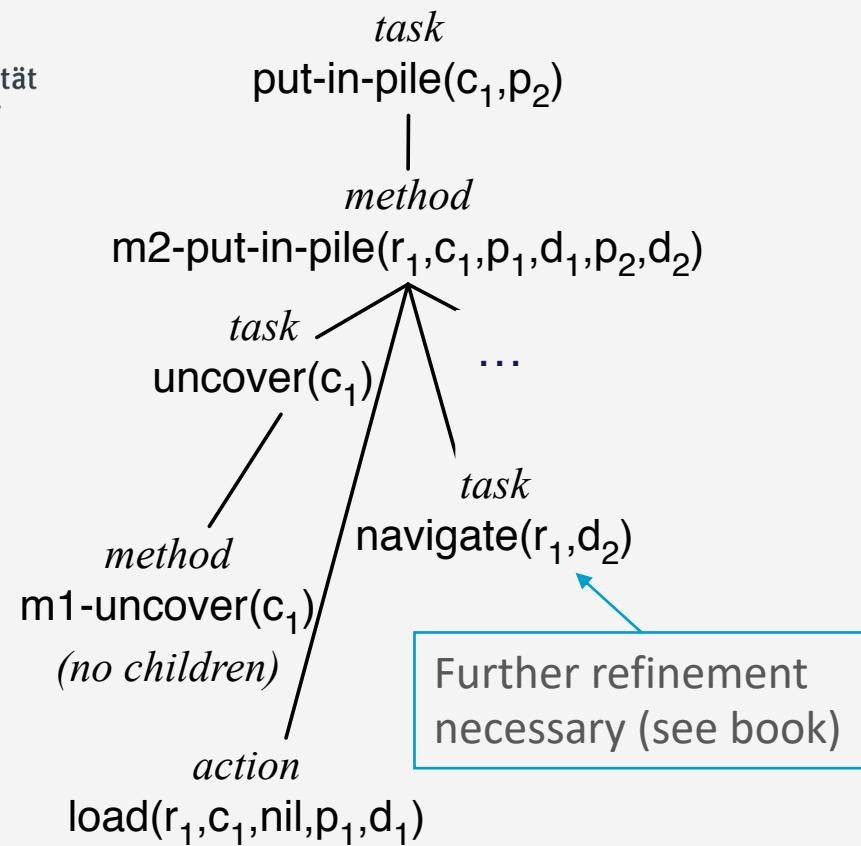
m2-put-in-pile(r, c, p, d, p', d')

...

body: if loc(r) ≠ d then navigate(r, d)
 uncover(c)
load(r, c, pos(c), p, d) **action**
 if loc(r) ≠ d' then navigate(r, d')
 unload(r, c, top(p'), p', d)



Example

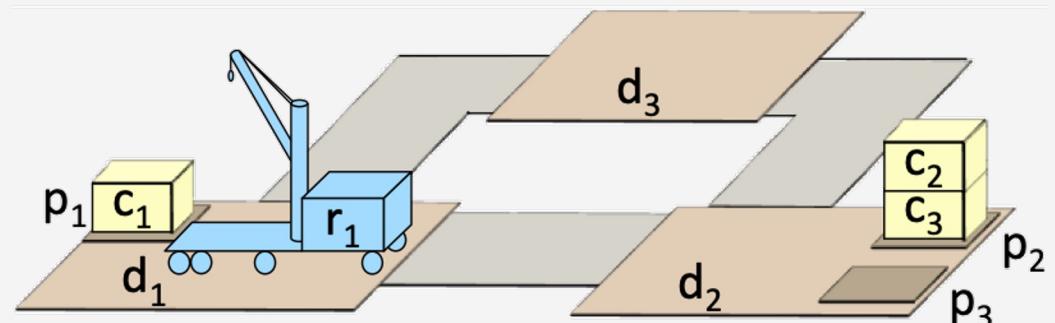


$r_1, c_1, p_1, d_1, p_2, d_2$

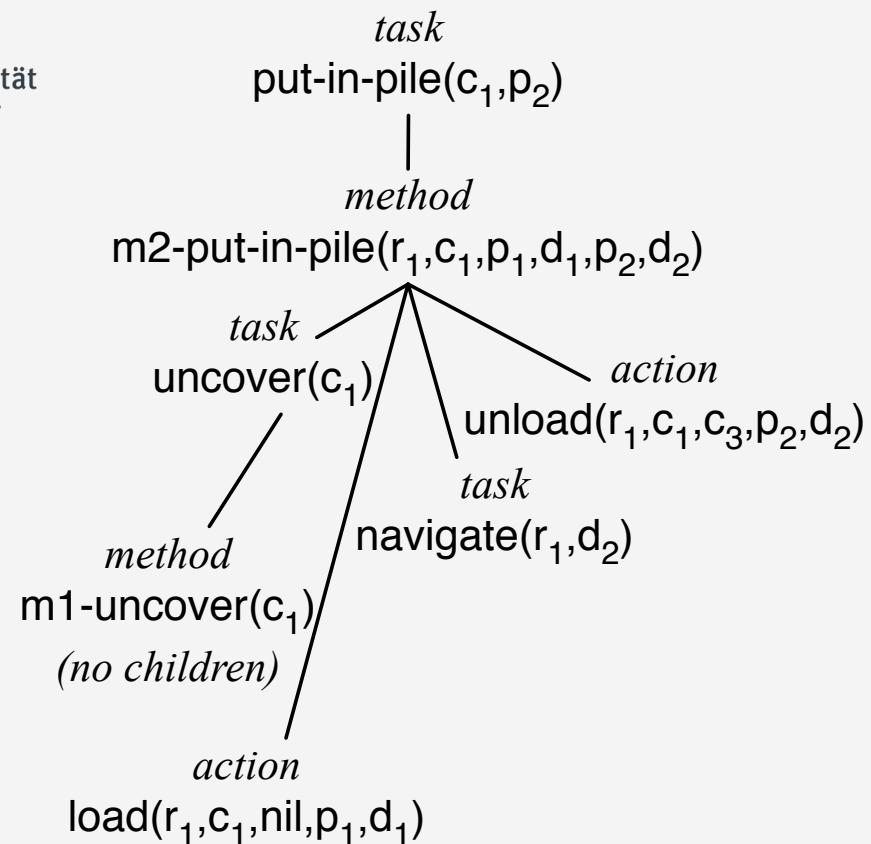
$\text{m2-put-in-pile}(r, c, p, d, p', d')$

...

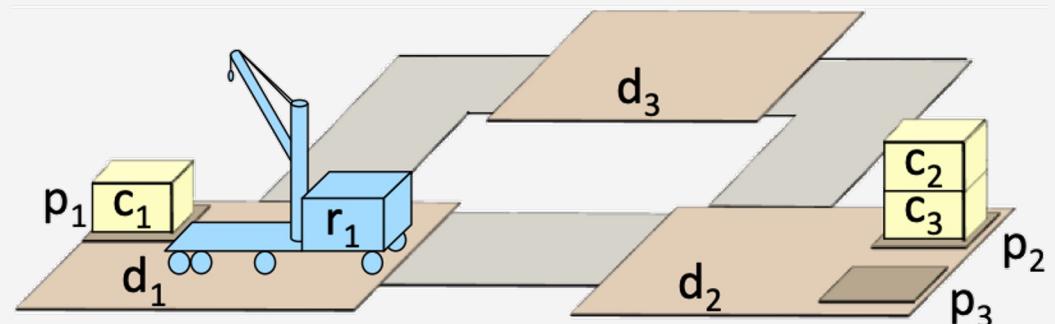
body: if $\text{loc}(r) \neq d$ then $\text{navigate}(r, d)$
 $\text{uncover}(c)$
 $\text{load}(r, c, \text{pos}(c), p, d)$ task
 if $\text{loc}(r) \neq d'$ then $\text{navigate}(r, d')$
 $\text{unload}(r, c, \text{top}(p'), p', d)$



Example



$r_1, c_1, p_1, d_1, p_2, d_2$
 $m2\text{-put-in-pile}(r, c, p, d, p', d')$
 ...
 body: if $\text{loc}(r) \neq d$ then $\text{navigate}(r, d)$
 $\text{uncover}(c)$
 $\text{load}(r, c, \text{pos}(c), p, d)$
 if $\text{loc}(r) \neq d'$ then $\text{navigate}(r, d')$
 $\text{unload}(r, c, \text{top}(p'), p', d)$ action



Heuristics For SeRPE

- *Ad hoc* approaches:
 - Domain-specific estimates
 - Statistical data on how well each method works
 - Try methods (or actions) in the order that they appear in \mathcal{M} (or \mathcal{A})
- Ideally, would want to implement using heuristic search (e.g., GBFS)
 - What heuristic function? Open problem
- SeRPE is a generalisation of HTN planning
 - In some cases, classical-planning heuristics can be used, in other cases they become intractable [Shivashankar *et al.*, ECAI-2016]

SeRPE ($\mathcal{M}, \mathcal{A}, s, \tau$)

$Candidates \leftarrow \text{Instances}(\mathcal{M}, \tau, s)$

if $Candidates = \emptyset$ **then**

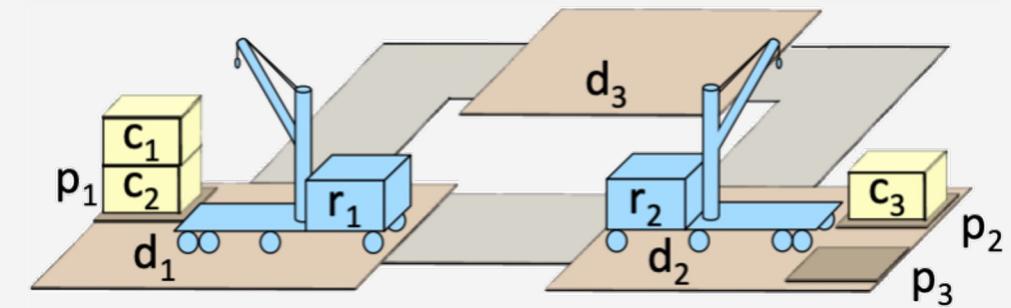
return failure

 nondeterministically choose $m \in Candidates$

return Progress-to-finish($\mathcal{M}, \mathcal{A}, s, \tau, m$)

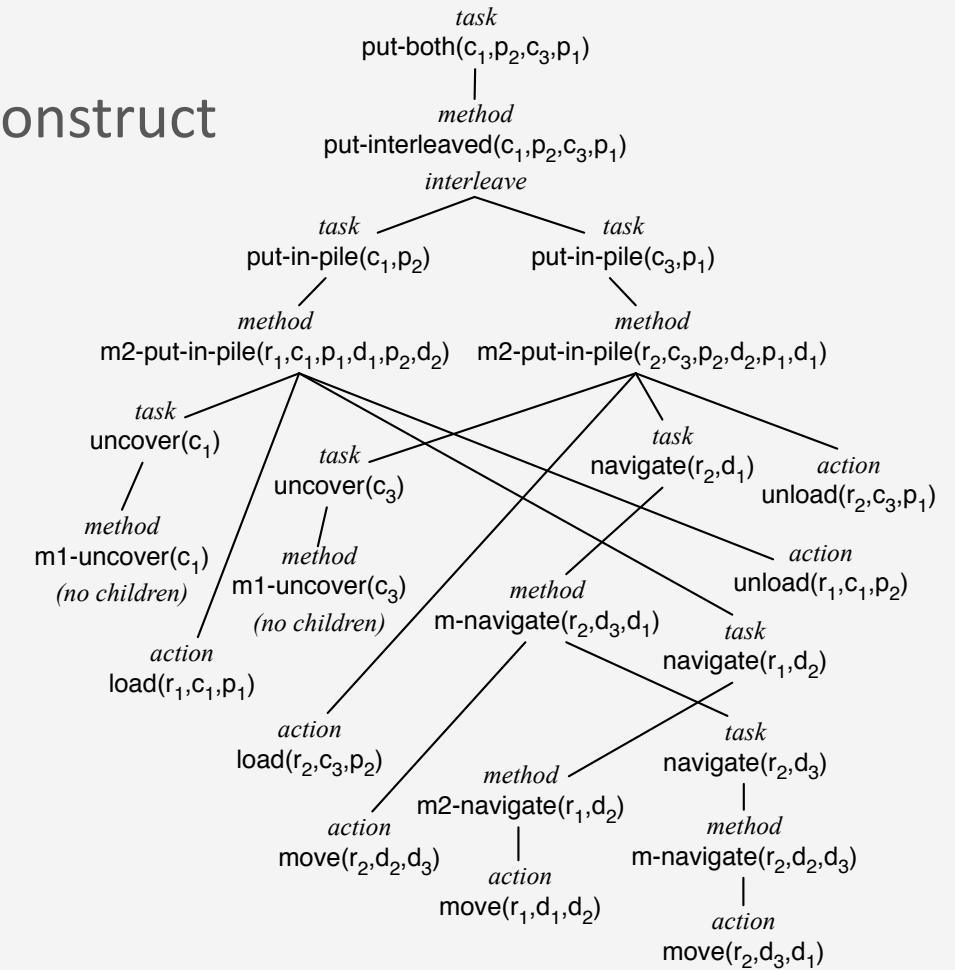
Interleaving

- Want to move c_1 to p_2 , using this plan
 - $\langle load(r_1, c_1, c_2, p_1, d_1), move(r_1, d_1, d_2), unload(r_1, c_1, p_3, nil, d_2) \rangle$
- ... and move c_3 to p_1 using this plan:
 - $\langle load(r_2, c_3, nil, p_2, d_2), move(r_2, d_2, d_3), move(r_2, d_3, d_1), unload(r_2, c_3, c_2, p_1, d_1) \rangle$
- For it to work, must interleave the plans
 - $\langle load(r_2, c_3, nil, p_2, d_2), move(r_2, d_2, d_3), load(r_1, c_1, c_2, p_1, d_1), move(r_1, d_1, d_2), unload(r_1, c_1, p_3, nil, d_2), move(r_2, d_3, d_1), unload(r_2, c_3, c_2, p_1, d_1) \rangle$



Interleaved Refinement Tree (IRT) Procedure

- SeRPE does not allow the “concurrent” programming construct
- Partial fix:
extend SeRPE to interleave plans for different tasks
- Details:
Section 3.3.2



Descriptive Action Models

- Predict the outcome of performing a command
 - Preconditions-and-effects representation
- Command
 - $take(r, o, l)$: robot r takes object o at location l
 - $put(r, o, l)$: r puts o at location l
 - $perceive(r, l)$: robot r perceives what objects are at location l
 - Can only perceive what is at its current location
 - If we knew this in advance, perception would not be necessary

- Action model

$take(r, o, l)$

pre: $cargo(r) = \text{nil}$, $\text{loc}(r) = l$, $\text{loc}(o) = l$

eff: $cargo(r) \leftarrow o$, $\text{loc}(o) \leftarrow r$

$put(r, o, l)$

pre: $\text{loc}(r) = l$, $\text{loc}(o) = r$

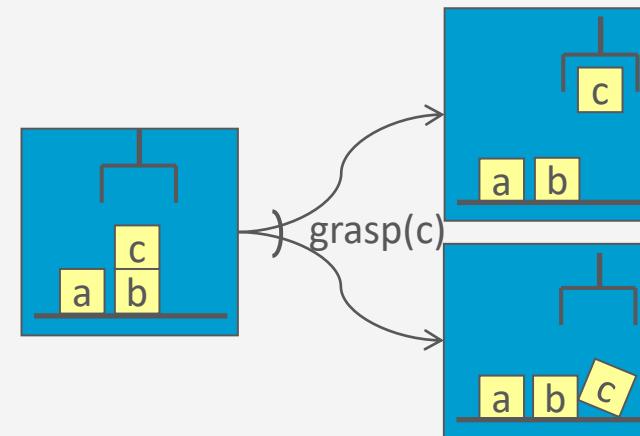
eff: $cargo(r) \leftarrow \text{nil}$, $\text{loc}(o) \leftarrow l$

$perceive(r, l)$

?

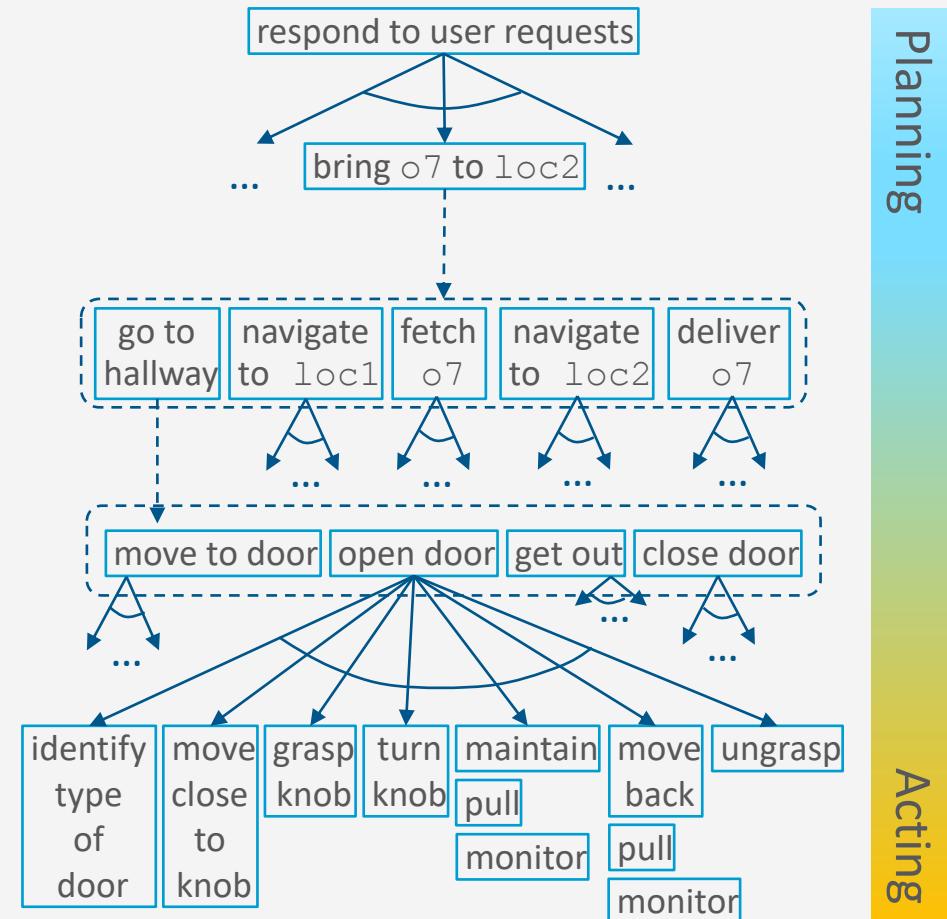
Limitation

- Most environments are inherently nondeterministic
 - Deterministic action models will not always make the right prediction
- Why use them?
 - Deterministic models \Rightarrow much simpler planning algorithms
 - Use when errors are infrequent and do not have severe consequences
 - Actor can fix the errors online



Planning/Acting at Different Levels

- Deterministic models may work better at some levels than others
- May want
 - RAE at some levels
 - RAE+planner at some levels
 - Planner at some levels
- In some cases, might want the planner to reason about nondeterministic outcomes
 - Later in lecture (Book: Ch. 5 + 6)
- Ongoing research on extending refinement planning to handle nondeterminism
[Patra *et al.*, AAAI-2019]



Summary

- Refinement planning (SeRPE)
 - Plan by simulating RAE on a single external task/event/goal
 - Deterministic actions
 - OK if we are confident of outcome, can recover if things go wrong
 - Interleaved plans (brief example)

Outline per the Book

3.1 *Representation*

- State variables, commands, refinement methods
- Example

3.2 *Acting*

- RAE (Refinement Acting Engine)
- Example
- Extensions

3.3 *Planning*

- Motivation and basic ideas
- Deterministic action models
- SeRPE (Sequential Refinement Planning Engine)

3.4 *Using Planning in Acting*

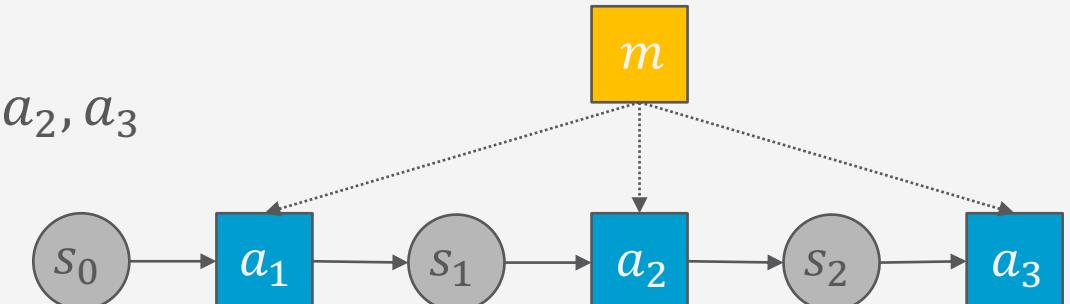
- Techniques
- Caveats

Acting and Refinement Planning

- Hierarchical acting with refinement planning
 - REAP: a RAE-like actor uses SeRPE-like planning at all levels
- Non-hierarchical actor with refinement planning
 - Refine-Lookahead, Refine-Lazy-Lookahead, Refine-Concurrent-Lookahead
 - Essentially the same as
 - Run-Lookahead, Run-Lazy-Lookahead, Run-Concurrent-Lookahead
 - But they call SeRPE instead of a classical planner for acting out a given task
 - Lookahead same as before
 - Receding horizon, sampling, subgoaling

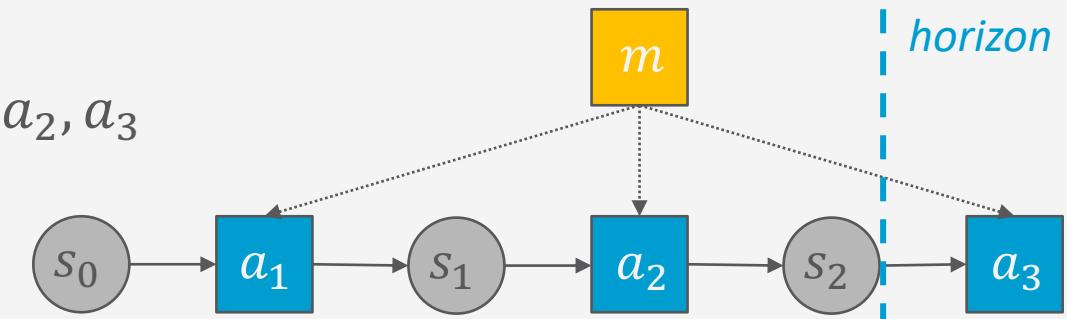
Caveats

- Start in state s_0 , want to accomplish task τ
 - Refinement method m with task τ , pre s_0 , body a_1, a_2, a_3
- Actor uses Refine-Lookahead
 - Lookahead = SeRPE, returns $\langle a_1, a_2, a_3 \rangle$
 - Actor performs a_1 , calls Lookahead again
 - No applicable method for τ in s_1 , SeRPE returns failure
- Fixes
 - When writing refinement methods, make them general enough to work in different states
 - In some cases, Lookahead might be able to fall back on classical planning until it finds something that matches a method
 - Keep snapshot of SeRPE's search tree at s_1 , resume there next time



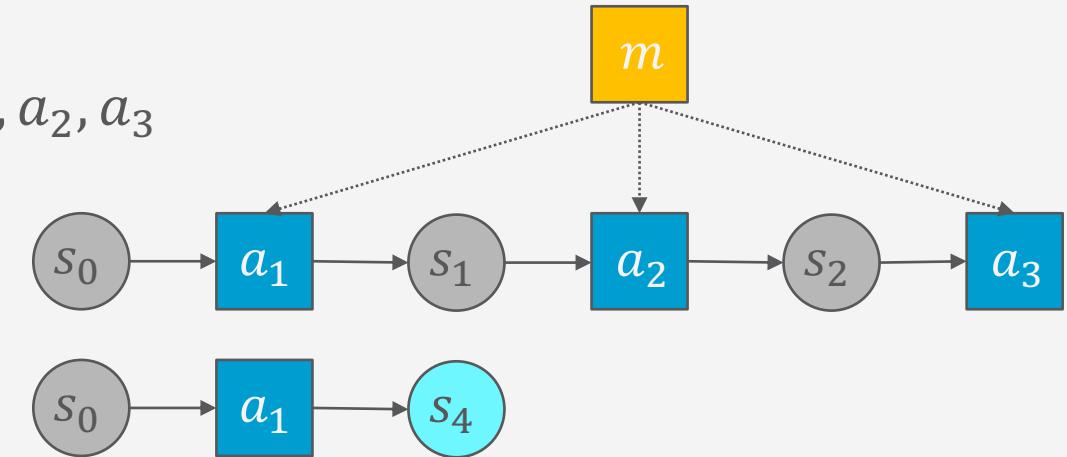
Caveats

- Start in state s_0 , want to accomplish task τ
 - Refinement method m with task τ , pre s_0 , body a_1, a_2, a_3
- Actor uses Refine-Lazy-Lookahead
 - Lookahead = SeRPE with receding horizon, returns $\langle a_1, a_2 \rangle$
 - Actor performs them, calls Lookahead again
 - No applicable method for τ in s_2 , SeRPE returns failure
- Fixes
 - Can use the same fixes on previous slide, with one modification
 - Keep snapshot of SeRPE's search tree **at the horizon**, resume next time it is called



Caveats

- Start in state s_0 , want to accomplish task τ
 - Refinement method m with task τ , pre: s_0 , body a_1, a_2, a_3
- Actor uses Refine-Lazy-Lookahead
 - Lookahead = SeRPE, returns $\langle a_1, a_2, a_3 \rangle$
 - While acting, unexpected event
 - Actor calls Lookahead again
 - No applicable method for τ in s_4 , SeRPE returns failure
- Fixes
 - Can use most of the fixes on last two slides, with this modification
 - Keep snapshot of SeRPE's search tree **after each action**
 - In example: restart it immediately after a_1 , using s_4 as current state
 - Also: make **recovery methods** for unexpected states



Summary

- Acting and planning
 - Lookahead: search part of the search space, return a partial solution
 - Refine-Lookahead, Refine-Lazy-Lookahead, Refine-Concurrent-Lookahead
 - Like Run-Lookahead, Run-Lazy-Lookahead, Run-Concurrent-Lookahead, but call SeRPE
 - Caveats
 - Current state may not be what we expect
 - Possible ways to handle that

Outline per the Book

3.1 *Representation*

- State variables, commands, refinement methods
- Example

3.2 *Acting*

- RAE (Refinement Acting Engine)
- Example
- Extensions

3.3 *Planning*

- Motivation and basic ideas
- Deterministic action models
- SeRPE (Sequential Refinement Planning Engine)

3.4 *Using Planning in Acting*

- Techniques
- Caveats

⇒ Next: Planning and Acting with Temporal Models