

# **WG 2.2 and semantics of programming languages**

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IFIP WG 2.2 Meeting  
15–18 September 2014, Munich

# IFIP WG 2.2 pre-history

**1964: IFIP Working Conference, Baden-bei-Wien**

## FORMAL LANGUAGE DESCRIPTION LANGUAGES FOR COMPUTER PROGRAMMING

*Proceedings of the*  
IFIP Working Conference on  
Formal Language Description Languages



1966

# IFIP WG 2.2 history

**1962**

- ▶ TC 2 – *Software: Theory and Practice*

**1965(!)**

- ▶ WG 2.2 – *Formal Description of Programming Concepts*
  - chair: T. B. Steel, Jr.

**1967**

- ▶ first WG 2.2 meeting (Alghero, Italy)

**1984**

- ▶ (member)

# WG 2.2

## AIMS

- ▶ The [primary] aim of the Working Group is to ***explicate programming concepts*** through the development, examination and comparison of various ***formal models of these concepts***.

## SCOPE

- ▶ The Working Group will investigate ***formalisms and models*** which represent different approaches to ***formal specification of programming concepts***. [...]

wg22.labri.fr

# Own contributions

early 1980s

- ▶ **abstract semantic algebras**

late 80s – 90s

- ▶ **action semantics**

late 90s – 2000s

- ▶ **modular SOS**

2010s

- ▶ **component-based semantics**

# Further IFIP WGs

**WG 1.3** (member since 1994, chair 1998–2003)

- ▶ ***Foundations of system specification***
- ▶ **CoFI**: Common Framework Initiative for algebraic specification and development of software
  - **CASL**

# Further IFIP WGs

**WG 2.11** (member since 2013)

▶ ***Program generation***

**SCOPE** [<https://wiki.hh.se/wg211/>]

- ▶ The working group covers the following research areas (and maybe others):
- ***programming language*** design, ***semantics*** and implementation
  - ***program synthesis***
  - ***type systems*** and type theory
  - [...]

# Recent WG 2.2 talks

- ▶ 2010 (Warsaw):
  - ***On bisimulation and modularity***
- ▶ 2011 (Paris):
  - ***PLANCOMPS – Programming Language Components and Specifications***
- ▶ 2012 (Amsterdam):
  - ***Component-based semantics***
- ▶ 2013 (Lisbon):
  - ***Editor support for formal specifications***



# This talk

## ***Component-based semantics***

- ▶ brief recap/overview
  - motivation
  - main ideas
- ▶ PLANCOMPS project
  - progress report
  - current and planned work

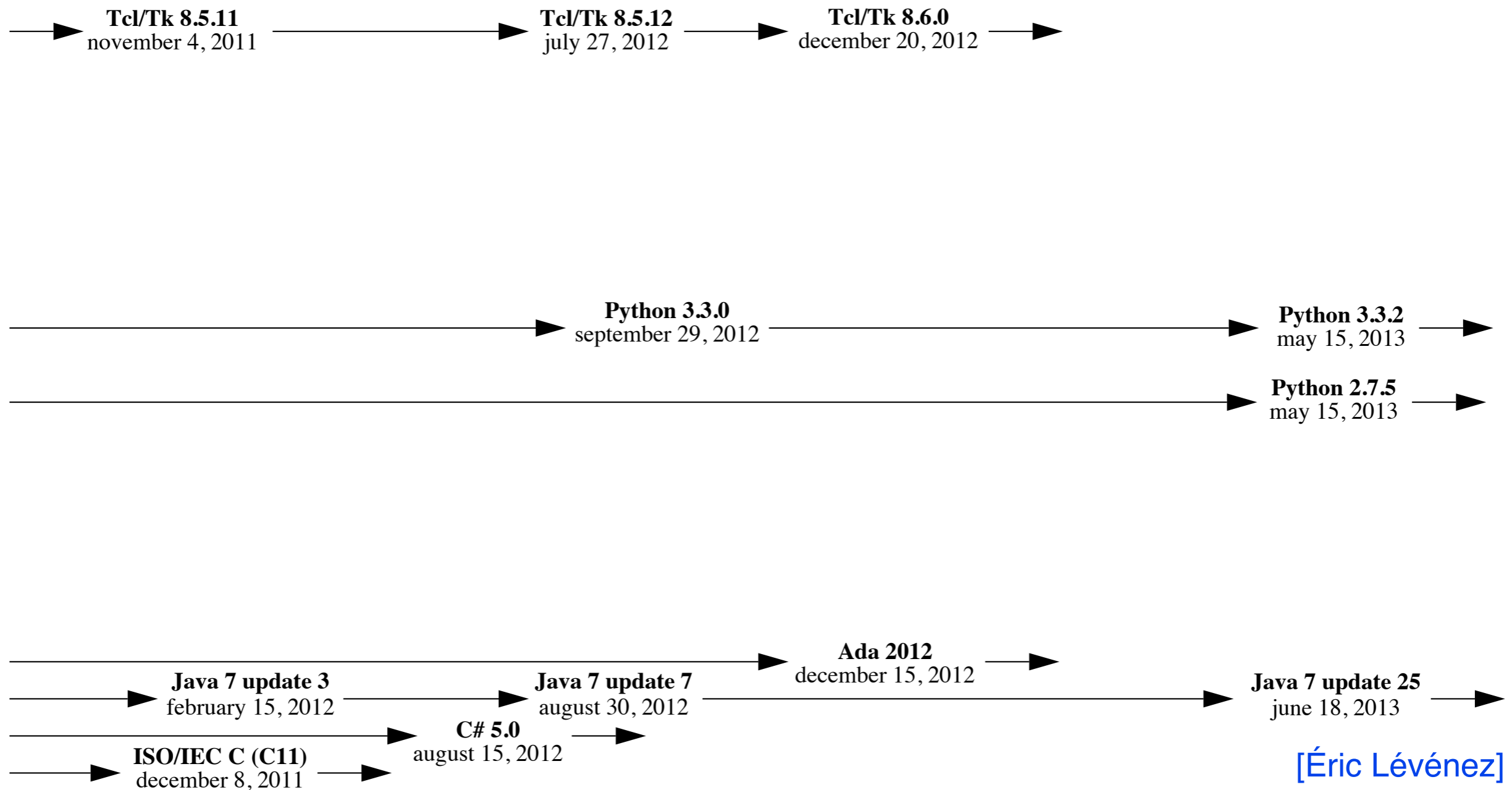
# **Component-based semantics**

# Programming languages

2012

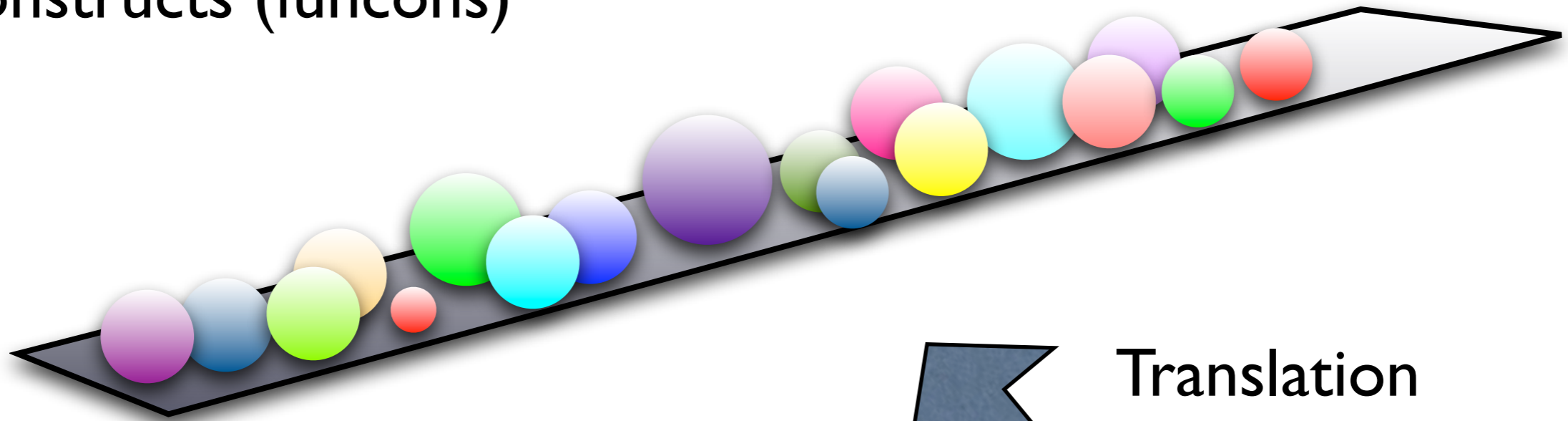
2013

## Evolution!



# Component-based semantics

Fundamental programming constructs (funcons)



Components-off-the-shelf (digital library)



Translation (reduction)

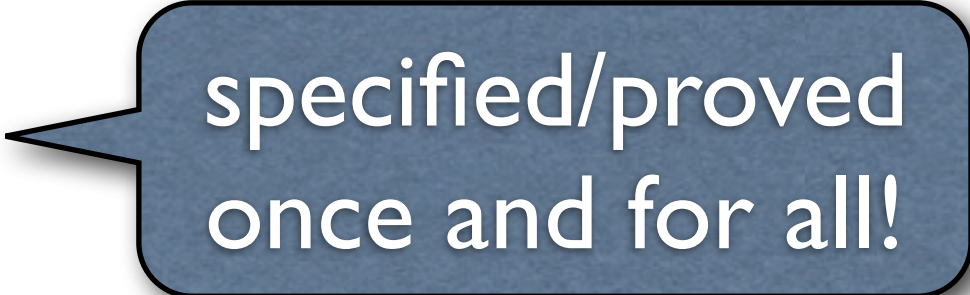


Evolving languages

# Reusable components

## Fundamental constructs (funcons)

- ▶ correspond to programming constructs
  - **directly** (**if-true**), or
  - **special case** (**apply**), or
  - **implicit** (**bound-value**)
- ▶ and have (*when validated and released*)
  - **fixed** notation, and
  - **fixed** behaviour, and
  - **fixed** algebraic properties



specified/proved  
once and for all!

# Component reuse

## Language construct:

- ▶  $exp ::= exp \ ? \ exp \ : \ exp$

## Translation to funcons:

- ▶  $expr \llbracket E_1 \ ? \ E_2 \ : \ E_3 \rrbracket =$   
 $\text{if-true}(expr \llbracket E_1 \rrbracket, expr \llbracket E_2 \rrbracket, expr \llbracket E_3 \rrbracket)$

## For languages with non-Boolean tests:

- ▶  $expr \llbracket E_1 \ ? \ E_2 \ : \ E_3 \rrbracket =$   
 $\text{if-true}(\text{not}(\text{equal}(expr \llbracket E_1 \rrbracket, 0)),$   
 $expr \llbracket E_2 \rrbracket, expr \llbracket E_3 \rrbracket)$

# Component reuse

## Language construct:

- ▶  $stm ::= \mathbf{if}(exp) \text{ } stm \text{ } \mathbf{else} \text{ } stm$

## Translation to funcons:

- ▶  $comm[\mathbf{if}(E_1) S_2 \mathbf{else} S_3] =$   
 $\mathbf{if-true}(expr[E_1], comm[S_2], comm[S_3])$

## For languages with non-Boolean tests:

- ▶  $comm[\mathbf{if}(E_1) S_2 \mathbf{else} S_3] =$   
 $\mathbf{if-true}(\mathbf{not}(\mathbf{equal}(expr[E_1], 0)),$   
 $comm[S_2], comm[S_3])$

destructive  
change

# Component specification

## Notation

modular extension

**if-true**(boolean, comp( $T$ ), comp( $T$ )) : comp( $T$ )

## Static semantics

$$\frac{E : \text{boolean}, \quad X_1 : T, \quad X_2 : T}{\text{if-true}(E, X_1, X_2) : T}$$

specified  
once and  
for all!

## Dynamic semantics

**if-true**(true,  $X_1$ ,  $X_2$ )  $\rightarrow$   $X_1$

**if-true**(false,  $X_1$ ,  $X_2$ )  $\rightarrow$   $X_2$



# **PLANCOMPS project**

# PLANCOMPS – foundations

## FOSSACS'13:

- ▶ bisimilarity  
congruence  
format
- ▶ preservation by  
disjoint extension

## Modular Bisimulation Theory for Computations and Values

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**Abstract.** For structural operational semantics (SOS) of process algebras, various notions of bisimulation have been studied, together with rule formats ensuring that bisimilarity is a congruence. For programming languages, however, SOS generally involves auxiliary entities (e.g. stores) and computed values, and the standard bisimulation and rule formats are not directly applicable.

Here, we first introduce a notion of bisimulation based on the distinction between computations and values, with a corresponding liberal congruence format. We then provide metatheory for a modular variant of SOS (MSOS) which provides a systematic treatment of auxiliary entities. This is based on a higher order form of bisimulation, and we formulate an appropriate congruence format. Finally, we show how algebraic laws can be proved sound for bisimulation with reference only to the (M)SOS rules defining the programming constructs involved in them. Such laws remain sound for languages that involve further constructs.

# PLANCOMPS – foundations

SLE'13:

- ▶ addressing deficiency of disambiguation annotations in SDF, Rascal, Spoofax

## Safe Specification of Operator Precedence Rules

Ali Afroozeh<sup>1</sup>, Mark van den Brand<sup>3</sup>, Adrian Johnstone<sup>4</sup>, Elizabeth Scott<sup>4</sup>,  
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**Abstract.** In this paper we present an approach to specifying operator precedence based on declarative disambiguation constructs and an implementation mechanism based on grammar rewriting. We identify a problem with existing generalized context-free parsing and disambiguation technology: generating a correct parser for a language such as OCaml using declarative precedence specification is not possible without resorting to some manual grammar transformation. Our approach provides a fully declarative solution to operator precedence specification for context-free grammars, is independent of any parsing technology, and is safe in that it guarantees that the language of the resulting grammar will be the same as the language of the specification grammar. We evaluate our new approach by specifying the precedence rules from the OCaml reference manual against the highly ambiguous reference grammar and validate the output of our generated parser.

# PLANCOMPS – foundations

ESOP'14:

- ▶ refocusing  
small-step  
(M)SOS rules

## Deriving Pretty-Big-Step Semantics from Small-Step Semantics

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**Abstract.** Big-step semantics for languages with abrupt termination and/or divergence suffer from a serious duplication problem, addressed by the novel ‘pretty-big-step’ style presented by Charguéraud at ESOP’13. Such rules are less concise than corresponding small-step rules, but they have the same advantages as big-step rules for program correctness proofs. Here, we show how to automatically derive pretty-big-step rules directly from small-step rules by ‘refocusing’. This gives the best of both worlds: we only need to write the relatively concise small-step specifications, but our reasoning can be big-step as well as small-step. The use of strictness annotations to derive small-step congruence rules gives further conciseness.

# PLANCOMPS – mini case study

## MODULARITY'14:

- ▶ component-based semantics of Caml Light
- ▶ partially validated (by empirical testing)
- ▶ detailed introduction to the approach
- ▶ preliminary tool support

## Reusable Components of Semantic Specifications

Martin Churchill    Peter D. Mosses    Paolo Torrini

# PLANCOMPS – tool support

## *Preliminary tool chain:*

### ▶ **SPOOFAX**

- parsing programs (SDF3, disambiguation, AST creation)
- translating ASTs to funcon terms (SDF3, Stratego-1.2)
- browsing and editing component-based specifications (SDF3, NaBL, Stratego-1.2)

### ▶ **PROLOG**

- translating MSOS rules for funcons to PROLOG
- running funcon terms

# PLANCOMPS – current, planned

## ***Case studies:***

- ▶ C# (*started*), Java

## ***Improved tool support:***

- ▶ generating SDF3 and Prolog from CBS
- ▶ refocusing MSOS rules

## ***Digital library:***

- ▶ historic semantics documents (*Cliff Jones, Newcastle*)
  - *ALGOL 60 in various frameworks (scan, OCR to PDF)*
- ▶ browsing and searching an open-access repository

# Conclusion

## ***Component-based semantics***

- ▶ translating programs to funcons
- ▶ specifying funcons in (I-)MSOS

## ***PLANCOMPS project (2011-2015)***

- ▶ foundations
- ▶ case studies
- ▶ tool support
- ▶ digital library